

The Decision Problems for Properties of Transformation Groups*

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Abstract

In this paper the fundamental result about the decision problems for properties of FRT-groups (i. e. the groups each of which is isomorphic with a group generated by a finite number of recursive transformations) has been proved: Let p be any algebraic property for groups such that there is a FRT-group G_1 which has the property p , a FRT-group G_2 which has not the property p , and G_2 is not isomorphic with any subgroup of any FRT-group which has the property p . Then the problem of deciding, for any given group P generated by a finite number of recursive transformations, whether or not the group G isomorphic with P has the property p is unsolvable.

A sequence of important consequences may be obtained from the fundamental result.

Let G be any abstract group. If there exists a transformation group $P = (A, B, \dots, C)$ generated by a finite number of recursive transformations A, B, \dots, C , such that $G \cong P$, then we say that P is a representative of G as a group generated by a finite number of recursive transformations, or P is a FRT-representative of G ; say that G is a group which has a representative as a group generated by a finite number of recursive transformations, or G is a FRT-group, and write that $G = G_P$.

In this paper the following fundamental result about the decision problems for algebraic properties of FRT-groups is proved.

Theorem Let p be any algebraic property for groups (i. e. a property for groups which is invariable under isomorphisms) such that there is a FRT-group G_1 which has the property p , a FRT-group G_2 which has not the property p , and G_2

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is not isomorphic with any subgroup of and FRT-group which has the property p . Then the problem of deciding, for any given FRT-representative P , whether or not G_P has the property p is unsolvable.

From this the following important consequences follows.

Corollary Let p be any algebraic property for groups which is hereditary and non-trivial with respect to FRT-groups (i. e. there is a FRT-group G_1 which has the property and a FRT-group G_2 which has not the property). Then the problem of deciding, for any given FRT-representative P , whether or not G_P has the property p is unsolvable.

Thus the following decision problems are all unsolvable: deciding, for any given FRT-representative P , whether or not G_P is

- (1) cyclic, (2) finite, (3) free, (4) commutative, (5) solvable,
(6) the group of units, (7) torsion-free group, and so on.

To prove the theorem we need a lemma.

Lemma Let \mathcal{P} be the set of all FRT-representatives, \mathcal{T} be the set of all recursive transformations on N . Then there exists a mapping τ from $\mathcal{P} \times \mathcal{T}$ into \mathcal{P} such that, for any $(P, T) \in \mathcal{P} \times \mathcal{T}$,

$$\tau(P, T) \simeq \begin{cases} \{e\}, & \text{if } T = I, \\ P, & \text{if } T \neq I, \end{cases} \quad (1)$$

where I is the identity transformation.

Proof For any recursive transformation S on N define transformations S_2 and S_3 on $N \times N \times N$ as follows:

$$S_2: (i, i, k) \rightarrow (i, iS, k); (i, iS, k) \rightarrow (i, i, k); (i, j, k) \rightarrow (i, j, k), \text{ if } j \neq i, iS.$$

$$S_3: (i, i, k) \rightarrow (i, i, kS); (i, j, k) \rightarrow (i, j, k), \text{ if } j \neq i.$$

For any recursive transformations S and T on N , define transformations S'_T and S_T^* on $N \times N \times N$ as follows:

$$S'_T = S_3 T_2 S_3^{-1} T_2.$$

$$S_T^*: (i, j, k) \rightarrow (i, j, k) S'_T, \text{ if } i = j, \text{ or } i \neq j \text{ and } (i, j) \in T;$$

$$(i, j, k) \rightarrow (i, j, k S^2) S'_T, \text{ if } i \neq j \text{ and } (i, j) \in T.$$

It is easy to see that S_2, S_3, S'_T, S_T^* are all recursive, and that

$$S'_T: (i, i, k) \rightarrow (i, i, k) \text{ if } (i, i) \in T; (i, i, k) \rightarrow (i, i, kS) \text{ if } (i, i) \notin T;$$

$$(i, j, k) \rightarrow (i, j, k S^{-1}) \text{ if } i \neq j \text{ and } (i, j) \in T; (i, j, k) \rightarrow (i, j, k) \text{ if } i \neq j \text{ and } (i, j) \notin T.$$

Therefore S_T^* : $(i, i, k) \rightarrow (i, i, k) \text{ if } (i, i) \in T; (i, i, k) \rightarrow (i, i, kS) \text{ if } (i, i) \notin T;$

$$(i, j, k) \rightarrow (i, j, kS), \text{ if } i \neq j \text{ and } (i, j) \in T;$$

$$(i, j, k) \rightarrow (i, j, k), \text{ if } i \neq j \text{ and } (i, j) \notin T.$$

Let f be any recursive correspondence between $N \times N \times N$ onto N , and write the image of $(i, j, k) \in N \times N \times N$ under f to be $\langle i, j, k \rangle \in N$. For any recursive transformations S and T on N define a transformation S_T on N as follows:

$$S_T = f^{-1} S_T^* f, \quad (2)$$

thus S_T is recursive,

For any FRT-representative $P = (A, B, \dots, C)$ and any recursive transformation T on N , form FRT-representative P_T as:

$P_T = (A_T, B_T, \dots, C_T)$. Let $\tau: \mathcal{P} \times \mathcal{T} \rightarrow \mathcal{P}$ be given by $\tau(P, T) = P_T$, we prove that τ satisfies the property (1).

If $T = I$, then, for any $i, j \in N$, $(i, j) \in T$ iff $i = j$. Hence, by (2), $A_T = B_T = \dots = C_T = I$, and therefore $P_T \simeq \{e\}$.

If $T \neq I$, then let $\varphi: P \rightarrow P_T$ satisfy, for any word $XY \dots Z \in P$ ($X, Y, \dots, Z \in \{A, A^{-1}, B, B^{-1}, \dots, C, C^{-1}\}$),

$$\varphi: XY \dots Z \rightarrow X_T Y_T \dots Z_T,$$

it is clear that φ is a mapping from P onto P_T , and $\langle i, j, k \rangle X_T Y_T \dots Z_T$

$$\begin{aligned} &= \begin{cases} \langle i, j, k \rangle Y_T \dots Z_T, & \text{if } i = j \text{ and } (i, j) \in T, \text{ or } i \neq j \text{ and } (i, j) \notin T \\ \langle i, j, kX \rangle Y_T \dots Z_T, & \text{if } i = j \text{ and } (i, j) \notin T, \text{ or } i \neq j \text{ and } (i, j) \in T \end{cases} \\ &= \begin{cases} \langle i, j, k \rangle, & \text{if } i = j \text{ and } (i, j) \in T, \text{ or } i \neq j \text{ and } (i, j) \notin T \\ \langle i, j, k(XY \dots Z) \rangle, & \text{if } i = j \text{ and } (i, j) \notin T, \text{ or } i \neq j \text{ and } (i, j) \in T \end{cases} \\ &= \langle i, j, k \rangle (XY \dots Z)_T. \end{aligned}$$

Hence P is homomorphic with P_T under φ .

Now, if $A = B = \dots = C = I$, then from $P \sim P_T$ it follows that $P \simeq P_T (\simeq \{e\})$. Otherwise, if $XY \dots Z \neq I$, then there is $k \in N$ such that $k \neq k(XY \dots Z)$.

Since $T \neq I$, hence there are m and n such that $m \neq n$ and $(m, n) \in T$. Thus $\langle m, n, k \rangle X_T Y_T \dots Z_T = \langle m, n, k(XY \dots Z) \rangle \neq \langle m, n, k \rangle$.

So that $X_T Y_T \dots Z_T \neq I$. Therefore we have $P \simeq P_T$ again. ■

Now we prove the theorem.

Proof of theorem Let us, for any recursive transformations S and T on N , write the recursive transformation in $\{(a, b) \mid (a = 2x \& b = 2y \& (x, y) \in S) \vee (a = 2x + 1 \& b = 2y + 1 \& (x, y) \in T)\}$ on N as $\begin{pmatrix} S \\ T \end{pmatrix}$, for any FRT-representative $P = (R_1, R_2, \dots, R_m)$ and $Q = (S_1, S_2, \dots, S_n)$, write the FRT-representative

$$\left(\begin{pmatrix} R_1 \\ I \end{pmatrix}, \begin{pmatrix} R_2 \\ I \end{pmatrix}, \dots, \begin{pmatrix} R_m \\ I \end{pmatrix}, \begin{pmatrix} I \\ S_1 \end{pmatrix}, \begin{pmatrix} I \\ S_2 \end{pmatrix}, \dots, \begin{pmatrix} I \\ S_n \end{pmatrix} \right)$$

as P^*Q .

Let the FRT-representatives of G_1 and G_2 be P_1 and P_2 respectively, and let P_0 be any FRT-representative which has an unsolvable word problem (for existence of P_0 see [1]). Form FRT-representative Q :

$$Q = P_1^* P_0.$$

Thus Q has an unsolvable word problem. For any word w of Q (which is obviously a recursive transformation on N), form FRT-representative $P(w)$:

$$P(w) = P_1^* Q_w.$$

Thus, if $w = I$, then $Q_w \simeq \{e\}$, $P(w) \simeq P_1$ has property p . If $w \neq I$, then $P(w) \simeq P_1^* Q$ (for $Q_w \simeq Q$). Hence P_2 is isomorphic with a subgroup of $P(w)$. By hypotheses, $P(w)$ has not property p .

Now, if the decision problem mentioned in the Theorem is solvable, the problem of deciding, for any $w \in Q$, whether or not $w = I$ is solvable, too. Hence from the fact that Q has an unsolvable word problem, our theorem follows. ■

Reference

- [1] Lin Yucai, Relation Word Problems.