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**Subcase 2'.2.6** A section of  $C_1$  is  $v_1, \dots, v_3, \dots, v_4$ , where the vertices not written are all vertices of  $S_1$ . Then the sequence of vertices

$$w, v_0, v_1, \dots, v_3, \dots, v_4, v_2, u, \dots, v_5, \dots, w$$

is a Hamiltonian cycle  $C$  of  $T(G)$ .

By above discussion, it is easy to see that  $T(G)$  has a Hamiltonian cycle when  $G$  has  $n$  cut vertices, and so the theorem follows.  $\square$

**Corollary 2.3** For any tree  $G$  of order not less than 2, the sufficient and necessary condition of  $T(G) \in H$  is that  $G$  is a path.

**Proof** By Theorem 1.3,  $T(G) \in P$ , and by Theorem 2.2, it is easy to see that the conclusion is true.  $\square$

## References

- [1] J.A.Bondy and U.S.R.Murty, *Graph Theory with Applications*, The Macmillan Press Ltd., 1976.
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## 关于完美全图的 Hamilton- 性

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### 摘 要

设  $G(V, E)$  是一个简单图, 而  $V(T(G)) = V(G) \cup E(G)$ ,  $E(T(G)) = \{yz | y, z \text{ 相邻或相关}, y, z \in V(G) \cup E(G)\}$ . 则称  $T(G)$  为  $G(V, E)$  的全图; 若对  $G$  的每一导出子图  $H$ , 有  $\chi(H) = \omega(H)$ , 则称  $G$  是完美的. 其中  $\chi(H)$ ,  $\omega(H)$  分别表示  $H$  的色数和团数. 本文给出了完美全图是 Hamilton 图的充分必要条件.

# On the Hamiltonity of Perfect Total Graphs \*

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**Abstract** Let  $G(V, E)$  be a simple graph, and  $V(T(G)) = V(G) \cup E(G)$ ,  $E(T(G)) = \{yz|y \text{ is adjacent or incident to } z, yz \in V(G) \cup E(G)\}$ . Let  $\chi(H)$  and  $\omega(H)$  be the chromatic number and the number of cliques of  $H$ , respectively. Then  $T(G)$  is called a total graph of  $G$ .  $G$  is called perfect if  $\chi(H) = \omega(H)$  for each induced subgraph  $H$  of  $G$ . We give a necessary and sufficient condition for a perfect total graph to be hamiltonian.

**Keywords** perfect total graph, Hamiltonity.

**Classification** AMS(1991)05C15/CCL O157.5

## 1 Introduction

**Definition 1.1** Let  $G(V, E)$  be a simple graph and

$$V(T(G)) = V(G) \cup E(G),$$

$$E(T(G)) = \{yz|y \text{ is adjacent or incident to } z, yz \in V(G) \cup E(G)\}.$$

Then  $T(G)$  is called a total graph of  $G$ ,  $E(G)$  is called derivative vertex set.

**Definition 1.2**<sup>[1]</sup> Let  $\chi(G)$  and  $\omega(G)$  be the chromatic number and the number of cliques of  $G$ , respectively;  $G[S]$  be an induced subgraph of  $G$  for  $S \subseteq V(G)$ ,  $G$  is called perfect if  $\chi(G[S]) = \omega(G[S])$  for any  $S \subseteq V(G)$ , and it is simply represented by  $G \in F$ .

**Theorem 1.3**<sup>[2]</sup> For any simple graph  $G$ ,  $T(G) \in F$ , if and only if each block of  $G$  includes at most three vertices.  $G(V, E)$  is called a hamiltonian graph or we say  $G(V, E)$  have hamiltonicity if graph  $G(V, E)$  has hamiltonian cycle; it is simply represented by  $G - H$ .

We refer to [1],[2] for notations and terminologies not explained here. The graphs that we consider here are simple.

## 2 The Main Results

**Lemma 2.1** For total graph  $K$  of order  $p$ ,

$$T(K_2), T(K_3) \in H.$$

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The conclusion is trivial. The proof is omitted here.

**Theorem 2.2** For any graph  $G, T(G) \in P$ , the necessary and sufficient condition of  $T(G) \in H$  is that any cut vertex of  $G$  is incident to at most two  $K_2$ -blocks.

**Proof** First we prove the necessary condition.

We will prove it by contradiction. If  $v_0 \in V(G)$  and it is a cut vertex of  $G$ . Then suppose  $v_0$  is incident to three  $K_2$ -blocks  $G[\{v_0, v_i\}] (i = 1, 2, 3)$ , where  $v_i \neq v_j (i \neq j, i, j = 1, 2, 3)$ . Since  $T(G) \in H$ , it is easy to see whether  $v_1, v_2, v_3$  are cut vertices of  $G$  or not, there will be the following conclusion: There are just two internally-disjoint paths between  $v_0$  and  $v_i (i = 1, 2, 3)$ .

Suppose that  $C$  is a Hamiltonian cycle of  $T(G)$ , and in cycle  $C$ .

$$d_C(v_0) = d_C(v_i) = 2 (i = 1, 2, 3).$$

Because  $v_0$  is a cut vertex of  $G$ , by Lemma 1 and  $v_1v_2, v_2v_3, v_3v_1 \in E(T(G))$ , we get  $v_1, v_2$  and  $v_3$  are all incident to  $v_0$  in cycle  $C$ . But by hypothesis  $v_i \neq v_j (i \neq j, i, j = 1, 2, 3)$ , this is impossible.

We can similarly prove that it is impossible for  $v_0$  to be incident to more than three  $K_2$ -blocks. So the necessary is true.

Now we prove the sufficient condition by induction on the number of cut vertices  $n$  of  $G$ .

When  $n = 0$ , by Theorem 1.3,  $G$  is  $K_2$  or  $K_3$ . From Lemma 2.1 the conclusion is true.

When  $n = 1$ , suppose  $v_0$  is the cut vertex. By Theorem 1.3 there are the following three cases.

**Case 1** There is no  $K_2$ -block in  $G$ .

From theorem 1.3, we know each block of  $G$  is  $K_3$  at the cut vertex  $v_0$ , the vertices of  $K_3$ -block are denoted (according to blocks) as  $1, 2, 3, 4, \dots, 2k-1, 2k$ ; where  $k$  is the number of blocks of  $K_3$  in graph  $G$ . The derivative vertices are represented by  $0 \cdot i (i = 1, 2, \dots, 2k)$  and  $i \cdot (i+1) (i = 1, 2, \dots, 2k-1)$ . Where  $0 \cdot i$  represents those from the edges between the vertices  $v_0$  and  $i (i = 1, 2, \dots, 2k)$  in  $T(G)$ ;  $1 \cdot (i+1)$  represents those from the edges between vertices  $v_i$  and  $i+1 (i = 1, 3, \dots, 2k-1)$  in  $T(G)$ . Then the sequence of vertices

$$v_0, 0 \cdot 1, 1, 1 \cdot 2, 2, 0 \cdot 2, 0 \cdot 3, 3, 3 \cdot 4, 4, 0 \cdot 4, \dots, 0 \cdot (2k-1), 2k-1,$$

$$(2k-1) \cdot (2k), 2k, 0 \cdot (2k), v_0$$

is a Hamiltonian cycle of  $T(G)$ .

**Case 2**  $G$  has just only one  $K_2$ -block.

First we give  $K_3$ -block the similar notation as case 1, the other vertices of  $K_2$ -block are represented by  $2k+1$ . They have the same meaning as case 1. Then the sequence of vertices

$$v_0, 2k+1, 0 \cdot (2k+1), 0 \cdot 1, 1, 1 \cdot 2, 2, \dots, 0 \cdot (2k-1), 2k-1,$$

$$(2k-1) \cdot (2k), 2k, 0 \cdot (2k), v_0$$

is a Hamiltonian cycle of  $T(G)$ .

**Case 3**  $G$  has two  $K_2$ -blocks.

Suppose there are  $kK_3$ -block in  $G$ .

**Subcase 3.1** When  $k = 0$ . The other two vertices of  $K_3$  are represented by 1,2. Then the sequence of vertices

$$v_0, 1, 0 \cdot 1, 0 \cdot 2, 2, v_0$$

is a Hamiltonian cycle of  $T(G)$ .

**Subcase 3.2** When  $k > 0$ .  $K_3$  has its vertices the same notation as case 1, the other vertices of two  $K_2$ -blocks are represented by  $2k+1, 2k+2$ , respaseively. Then the sequence of vertices

$$v_0, 2k+1, 0 \cdot (2k+1), 0 \cdot 1, 1, \dots, 2k, 0 \cdot (2k), 0 \cdot (2k+2), 2k+1, v_0$$

is a Hamiltonian cycle of  $T(G)$ .

By above discussion the conclusion is true for  $n = 1$ .

Suppose when the number of cut vertices is less than  $n$ , it is true that  $T(G) \in H$ . We now prove that  $T(G) \in H$  for the number  $n$  of cut vertices.

From the property of  $G$ , and since  $G$  is a limited graph,  $G$  must have a cut vertex  $v_0$  which is adjacent to at most two vertices, and let  $v_1$  be a cut vertex adjacent to it. We consider the following two cases.

**Case 1'**  $v_0v_1$  is a cut edge of  $G$ .

Let  $G_1$  be an induced subgraph of  $G$  with vertices of the branch including  $v_1$  of  $G - v_0$  and vertex  $v_0$ ;  $G_2$  be an induced subgraph of  $G$  with vertices of the branch incoluding  $v_0$  of  $G - v_1$  and vertex  $v_1$ . Then  $G_2$  has one cut vertex  $v_0$ ,  $G_1$  has  $n-1$  cut vertices. By the induction hypothesis,

$$T(G_1), T(G_2) \in H.$$

Let  $v_2$  represent the derivative vertex from edge  $v_0v_1$  of  $G$  in  $T(G)$ ,  $C_1$  the Hamiltonian cycle of  $T(G_i)$  ( $i = 1, 2$ ). Since

$$d_{T(G_1)}(v_0) = d_{T(G_2)}(v_1) = 2,$$

edges  $v_0v_1, v_0v_2$  are in the cycle  $C_1$ ,  $v_1v_0, v_1v_2$  are in the cycle  $C_2$ . We construct a cycle  $C$  based on  $C_1, C_2$  and let  $C$  satisfy

$$\begin{aligned} V(C) &= V(C_1) \cup V(C_2), \\ E(C) &= [E(C_1) \cup E(C_2)] \setminus \{v_0v_2, v_1v_2\}. \end{aligned}$$

Then the cycle  $C$  is a Hamiltonian cycle of  $T(G)$ .

**Case 2'**  $v_0v_1$  is not a cut edge of  $G$ .

By the conditions given in the theorem,  $G$  must have a vertex  $v_2$  such that  $v_0v_2, v_1v_2 \in E(G)$ . Otherwise,  $v_0v_1$  will be a cut edge of  $G$ . This is a contradiction.

**Subcase 2'.1** If  $v_2$  is not a cut vertex of  $G$ .

Let  $G_1$  represent the induced subgraph of  $G$  with the vertex  $v_0$  and vertices which belong to the branch of  $G - v_0$  including  $v_1$ ,  $G_2$  the induced subgraph of  $G$  with the vertex  $v_0$  and vertices of  $V(G) \setminus V(G - v_0)$ ;  $v_3, v_4, v_5$  the derivative vertices from edges  $v_0v_1, v_1v_2, v_2v_0$  in  $T(G)$ ; and

$$S_0 = \{v_0, v_1, v_2, v_3, v_4, v_5\};$$

$$S_1 = V(T(G_1)) \setminus S_0;$$

$$S_2 = V(T(G_2)) \setminus \{v_0\}.$$

Then  $G_2$  has at most one cut vertex  $G_1$  has  $n-1$  cut vertices. By the induction hypothesis,  $T(G_i) \in H(i=1, 2)$ . From  $n=0$  or  $n=1$ , we can suppose that  $v_0, u, \dots, w, v_0$  is a Hamiltonian cycle  $C_2$  of  $T(G_2)$ , where  $u, w \in V(T(G_2))$ , the other vertices (not written) are all vertices of  $S_2 \setminus \{u, w\}$ , and  $v_5u, wv_3 \in E(T(G))$ .

Let  $C_1$  represent a hamiltonian cycle of  $T(G_1)$ . We consider that

$$d_{T(G_1)}(v_1) = 4(i=0, 2, 5)$$

and the vertices  $v_0, v_2, v_5$  are not adjacent to any one of  $S_1$  in  $T(G_1)$ , then we have the following six subcases.

**Subcase 2'.1.1** A section of  $C_1$  is  $v_3, \dots, v_1$ , where the vertices not written are all vertices of  $S_1$ . Then the sequence of vertices  $w, v_3, \dots, v_1, v_4, v_2, v_5, v_0, u, \dots, w$  is a Hamiltonian cycle  $C$  of  $T(G)$ .

**Subcase 2'.1.2** A section of  $C_1$  is  $v_3, \dots, v_4$ , where the vertices not written are all vertices of  $S_1$ . Then the sequence of vertices  $w, v_3, \dots, v_4, v_1, v_2, v_5, v_0, u, \dots, w$  is a Hamiltonian cycle  $C$  of  $T(G)$ .

**Subcase 2'.1.3** A section of  $C_1$  is  $v_1, \dots, v_4$ , where the vertices not written are all vertices of  $S_1$ . Then the sequence of vertices  $w, v_3, v_1, \dots, v_4, v_2, v_5, v_0, u, \dots, w$  is a Hamiltonian cycle  $C$  of  $T(G)$ .

**Subcase 2'.1.4** A section of  $C_1$  is  $v_3, \dots, v_1, \dots, v_4$ , where the vertices not written are all vertices of  $S_1$ . Then the sequence of vertices  $w, v_3, \dots, v_1, \dots, v_4, v_2, v_5, v_0, u, \dots, w$  is a Hamiltonian cycle  $C$  of  $T(G)$ .

**Subcase 2'.1.5** A section of  $C_1$  is  $v_3, \dots, v_4, \dots, v_1$ , where the vertices not written are all vertices of  $S_1$ . Then the sequence of vertices  $w, v_3, \dots, v_4, \dots, v_1, v_2, v_5, v_0, u, \dots, w$  is a Hamiltonian cycle  $C$  of  $T(G)$ .

**Subcase 2'.1.6** A section of  $C_1$  is  $v_1, \dots, v_3, \dots, v_4$ , where the vertices not written are all vertices of  $S_1$ . Then the sequence of vertices  $w, v_0, v_1, \dots, v_3, \dots, v_4, v_2, v_5, u, \dots, w$  is a Hamiltonian cycle  $C$  of  $T(G)$ .

By above discussions, it is easy to see that the conclusion is true when  $v_2$  is not a cut vertex of  $G$ .

**Subcase 2'.2** Each cut vertex of  $G$  is at least adjacent to two cut vertices.

As  $G$  is a limited graph. It's easy to prove by contradiction that  $G$  must have two adjacent cut vertices  $v_0, v_2$  which have only one common adjacent cut vertex  $v_1$ .

Let  $G_1$  represent the induced subgraph of  $G$  with vertices  $v_0, v_2$  and vertices of the branch of  $G - \{v_0, v_2\}$  including vertex  $v_1$ ,  $G_2$  the induced subgraph of  $G$  with vertex  $v_1$  and vertices of the branches of  $G - v_1$  including vertex  $v_0$  and  $v_2$ ;  $v_3, v_4, v_5$  have the same meaning as subcase 2.1, and

$$\begin{aligned} S_0 &= \{v_0, v_1, v_2, v_3, v_4\}; \\ S_i &= V(T(G_i)) \setminus (S_0 \cup \{v_5\}), \quad (i = 1, 2). \end{aligned}$$

Suppose  $u$  is a derivative vertex from one of edges incident to vertex  $v_2$ ;  $w$  is a derivative vertex from one of edges incident to vertex  $v_0$ ,  $u, w \in \{v_3, v_4, v_5\}$ ; and the sequence of vertices  $u, \dots, v_5, \dots, w$  represents a hamiltonian path including all vertices of  $S_2$  and  $v_5$  from  $u$  to  $w$ . Evidently, this path must exist.

Because  $G_2$  has two cut vertices  $v_0, v_2$  and  $G_1$  has  $n - 2$  cut vertices,  $T(G_1)$  must have a hamiltonian cycle  $C_1$  by the hypothesis.

We consider  $C_1$  as the following six subcases similarly to subcase 2'.1.

**Subcase 2'.2.1** A section of  $C_1$  is  $v_1, \dots, v_3$ , where the vertices not written are all vertices of  $S_1$ . Then the sequence of vertices

$$w, v_0, v_3, \dots, v_1, v_4, v_2, u, \dots, v_5, \dots, w$$

is a Hamiltonian cycle  $C$  of  $T(G)$ .

**Subcase 2'.2.2** A section of  $C_1$  is  $v_1, \dots, v_4$ , where the vertices not written are all vertices of  $S_1$ . Then the sequence of vertices

$$w, v_0, v_3, v_1, \dots, v_4, v_2, u, \dots, v_5, \dots, w$$

is a Hamiltonian cycle  $C$  of  $T(G)$ .

**Subcase 2'.2.3** A section of  $C_1$  is  $v_3, \dots, v_4$ , where the vertices not written are all vertices of  $S_1$ . Then the sequence of vertices

$$w, v_0, v_1, v_3, \dots, v_4, v_2, u, \dots, v_5, \dots, w$$

is a Hamiltonian cycle  $C$  of  $T(G)$ .

**Subcase 2'.2.4** A section of  $C_1$  is  $v_3, \dots, v_1, \dots, v_4$ , where the vertices not written are all vertices of  $S_1$ . Then the sequence of vertices

$$w, v_0, v_3, \dots, v_1, \dots, v_4, v_2, u, \dots, v_5, \dots, w$$

is a Hamiltonian cycle  $C$  of  $T(G)$ .

**Subcase 2'.2.5** A section of  $C_1$  is  $v_3, \dots, v_4, \dots, v_1$ , where the vertices not written are all vertices of  $S_1$ . Then the sequence of vertices

$$w, v_0, v_3, \dots, v_4, \dots, v_1, v_2, u, \dots, v_5, \dots, w$$

is a Hamiltonian cycle  $C$  of  $T(G)$ .

**Subcase 2'.2.6** A section of  $C_1$  is  $v_1, \dots, v_3, \dots, v_4$ , where the vertices not written are all vertices of  $S_1$ . Then the sequence of vertices

$$w, v_0, v_1, \dots, v_3, \dots, v_4, v_2, u, \dots, v_5, \dots, w$$

is a Hamiltonian cycle  $C$  of  $T(G)$ .

By above discussion, it is easy to see that  $T(G)$  has a Hamiltonian cycle when  $G$  has  $n$  cut vertices, and so the theorem follows.  $\square$

**Corollary 2.3** For any tree  $G$  of order not less than 2, the sufficient and necessary condition of  $T(G) \in H$  is that  $G$  is a path.

**Proof** By Theorem 1.3,  $T(G) \in P$ , and by Theorem 2.2, it is easy to see that the conclusion is true.  $\square$

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