

Incorrect Results for E -Convex Functions and E -Convex Programming *

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Abstract: A class of functions and a sort of nonlinear programming called respectively E -convex functions and E -convex programming were presented and studied recently by Youness in [1]. In this paper, we point out the most results for E -convex functions and E -convex programming in [1] are not true by six counter examples.

Key words: Convex sets; convex functions; convex programming; generalized convexity.

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1. Definitions and some results in [1]

In this section, we recall the related definitions and results given in [1] which will be used in our study.

Definition 1 (Def. 2.1 in [1]) A set $M \subseteq R^n$ is said to be an E -convex set, if there exists a map $E : R^n \rightarrow R^n$ such that

$$\lambda E(x) + (1 - \lambda)E(y) \in M, \quad \forall x, y \in M, \forall \lambda \in [0, 1]. \quad (1)$$

Proposition 1 If set $M \subseteq R^n$ is an E -convex set, then $E(M) \subseteq M$.

Definition 2 (Def. 3.1 in [1]) A function $f : R^n \rightarrow R$ is said to be E -convex on a set $M \subseteq R^n$ if there is a map $E : R^n \rightarrow R^n$ such that M is an E -convex set and

$$f(\lambda E(x) + (1 - \lambda)E(y)) \leq \lambda f(E(x)) + (1 - \lambda)f(E(y)), \quad \forall x, y \in M, \lambda \in [0, 1]. \quad (2)$$

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Furthermore, if the inequality signs in formula (2) is strict for $E(x) \neq E(y)$ and $\lambda \in (0, 1)$, then f is called strictly E -convex.

Remark 1 The definition of strict E -convexity in [1] is not clear and definite, since the conditions for formula (2) being strict were not given.

Definition 3 (Def. 3.2 in [1]) Let $S \subseteq R^n \times R$ and $E : R^n \rightarrow R^n$. The set S is said to be E -convex if $(x, \alpha), (y, \beta) \in S$ imply

$$(\lambda E(x) + (1 - \lambda)E(y), \lambda\alpha + (1 - \lambda)\beta) \in S, \forall \lambda \in [0, 1]. \quad (3)$$

Definition 4 (See Section 4 in [1]) The nonlinear programming problem

$$(P) \quad \begin{array}{ll} \min & f(x), \\ \text{s.t.} & x \in M = \{x \in R^n : g_i(x) \leq 0, i \in I\}. \end{array}$$

is said to be an E -convex programming if there exists a map $E : R^n \rightarrow R^n$ such that the functions $f, g_i (i \in I) : R^n \rightarrow R$ all are E -convex functions on R^n .

Throughout this paper, Problem (P) is always assumed to be an E -convex programming.

Theorem 1 (Theorem 3.1 in [1]) A numerical function f defined on an E -convex set $M \subseteq R^n$ is E -convex on M iff its E -epigraph $E - e(f)$ is E -convex on $R^n \times R$, where

$$E - e(f) = \{(x, \alpha) : x \in M, \alpha \in R, f(E(x)) \leq \alpha\}.$$

Theorem 2 (Theorem 4.1 in [1]) The feasible set M of (P) is an E -convex set.

Theorem 3 (Theorem 4.2 in [1]) Assume that $E(M)$ is convex and \bar{x} is a solution of the following problem:

$$(P_E) \quad \min\{(f \circ E)(x) \mid x \in M\}.$$

Then $E(\bar{x})$ is a solution of Problem (P), where $(f \circ E)(x) = f(E(x))$.

Theorem 4 (Theorem 4.3 in [1]) Let $E(M)$ be a convex set. If $x^o = E(z^o) \in E(M)$ is a local minimum of Problem (P) on M , then x^o is a global minimum of Problem (P) on M .

Theorem 5 (Theorem 4.4 in [1]) Assume that $E(M)$ is a convex and f is strictly E -convex. Then, the solution of Problem (P_E) is unique.

Theorem 6 (Theorem 4.5 in [1]) Let $E(M)$ is a convex set and let $f \circ E, g_i \circ E (i \in I)$ all be differentiable on M . Assume that (x^*, y^*) is a solution of the following problem:

$$\begin{cases} \nabla_x[(f \circ E)(x^*) + y^*(g_i \circ E)(x^*)] = 0, \\ y^*(g_i \circ E)(x^*) = 0, (g_i \circ E)(x^*) \leq 0, y^* \geq 0, \forall i \in I. \end{cases} \quad (4)$$

Then, $E(x^*)$ is an optimal solution of Problem (P).

Theorem 7 (Theorem 4.6 in [1]) Let Ω be the set of optimal solution of (P). Then Ω is a convex set.

2. Counter examples for the above theorems

In this Section, we give six examples to show that the seven theorems given above are incorrect.

Example 1 A counter example for the necessity of Theorem 1.

Consider set M and maps $E, f : R \rightarrow R$ defined as

$$M = R = (-\infty, +\infty), \quad E(x) = -x^2, \quad f(x) = \begin{cases} 1, & \text{if } x > 0; \\ -x, & \text{if } x \leq 0. \end{cases}$$

Then M is an E -convex set, and f is an E -convex function on M since

$$f(\lambda E(x) + (1 - \lambda)E(y)) = \lambda f(E(x)) + (1 - \lambda)f(E(y)), \quad \forall x, y \in M, \forall \lambda \in [0, 1].$$

But the E -epigraph

$$E - e(f) = \{(x, \alpha) : x \in M, \alpha \in R, f(E(x)) = x^2 \leq \alpha\},$$

is not E -convex on R^2 , since for $(x, \alpha) = (1, 1), (y, \beta) = (2, 4) \in E - e(f)$ and $\lambda = \frac{1}{3}$, one has

$$(\lambda E(x) + (1 - \lambda)E(y), \lambda \alpha + (1 - \lambda)\beta) = (-3, 3), \quad (-3)^2 \not\leq 3,$$

that is $(\lambda E(x) + (1 - \lambda)E(y), \lambda \alpha + (1 - \lambda)\beta) = (-3, 3) \notin E - e(f)$.

Example 2 A counter example for the sufficiency of Theorem 1.

Let $M = [0, 1] \subset R, f(x) = -x^2, E(x) = \sqrt{x}, x \in M$.

Then M is an E -convex set, and

$$E - e(f) = \{(x, \alpha) : 0 \leq x \leq 1, \alpha \in R, f(E(x)) = -x \leq \alpha\}.$$

At first, we prove that the set $E - e(f)$ is E -convex on $R \times R$ as follows. Let $(x, \alpha), (y, \beta) \in E - e(f), \lambda \in [0, 1]$, then

$$-\sqrt{x} \leq -x \leq \alpha, \quad -\sqrt{y} \leq -y \leq \beta, \quad \lambda E(x) + (1 - \lambda)E(y) \in M,$$

$$(\lambda E(x) + (1 - \lambda)E(y), \lambda \alpha + (1 - \lambda)\beta) = (\lambda \sqrt{x} + (1 - \lambda)\sqrt{y}, \lambda \alpha + (1 - \lambda)\beta),$$

$$-\lambda \sqrt{x} - (1 - \lambda)\sqrt{y} \leq \lambda \alpha + (1 - \lambda)\beta.$$

Thus

$$(\lambda E(x) + (1 - \lambda)E(y), \lambda \alpha + (1 - \lambda)\beta) \in E - e(f).$$

Therefore set $E - e(f)$ is E -convex on $R \times R$ from Definition 3.

Secondly, one can conclude the given function f is not E -convex on the E -convex set M , since for $x = 0 \in M, y = 1 \in M, \lambda = \frac{1}{2}$, one gets

$$f(\lambda E(x) + (1 - \lambda)E(y)) = f\left(\frac{1}{2}\right) = -\frac{1}{4} > -\frac{1}{2} = \lambda f(E(x)) + (1 - \lambda)f(E(y)).$$

Example 3 A counter example for Theorem 2, i.e., Theorem 4.1 in [1].

Let maps $g, E : R \rightarrow R$ be given as

$$g(x) = \begin{cases} x, & \text{if } x < 0; \\ 1 - x, & \text{if } x \geq 0, \end{cases} \quad E(x) = |x|.$$

Then the function g is E -convex on R , see Fig. 1, since

$$g(\lambda E(x) + (1 - \lambda)E(y)) = \lambda g(E(x)) + (1 - \lambda)g(E(y)), \quad \forall x, y \in R, \lambda \in [0, 1].$$

But the feasible set

$$M = \{x : g(x) \leq 0\} = (-\infty, 0) \cup [1, +\infty)$$

is not E -convex from Proposition 1, since $E(M) = (0, +\infty) \not\subset M$.

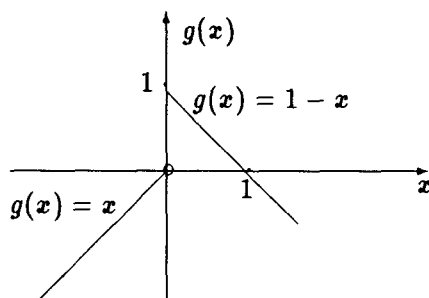


Fig.1 Example 3. A counter example for Theorem 2, i.e., Theorem 4.1 in [1]

Example 4 A counter example for Theorems 3,4, i.e., Theorems 4.2, 4.3 in [1].

Consider functions $f, g, E : R \rightarrow R$ defined as

$$f(x) = \begin{cases} -2, & \text{if } x \leq -1; \\ 2x, & \text{if } -1 < x < 0; \\ -x, & \text{if } 0 \leq x < 1; \\ -1, & \text{if } x \geq 1, \end{cases} \quad g(x) = \begin{cases} 1 + x, & \text{if } x < 0; \\ 1 - x, & \text{if } x \geq 0, \end{cases} \quad E(x) = |x|.$$

(i) Show that Thm. 3 is not true. See Fig. 2, one knows that

$$M = \{x : g(x) \leq 0\} = (-\infty, -1] \cup [1, +\infty), \quad E(R) = [0, +\infty), \quad E(M) = [1, +\infty),$$

$$(f \circ E)(x) = f(|x|) = \begin{cases} -|x|, & \text{if } 0 \leq |x| < 1; \\ -1, & \text{if } |x| \geq 1, \end{cases}$$

furthermore, f, g all are E -convex on R , $E(M)$ is a convex set .

It is clear that $\bar{x} = -2$ is a solution of (P_E) , see Fig. 2, but $E(\bar{x}) = 2$ is not a solution of Problem (P) since $f(2) = -1 > f(-2) = -2$.

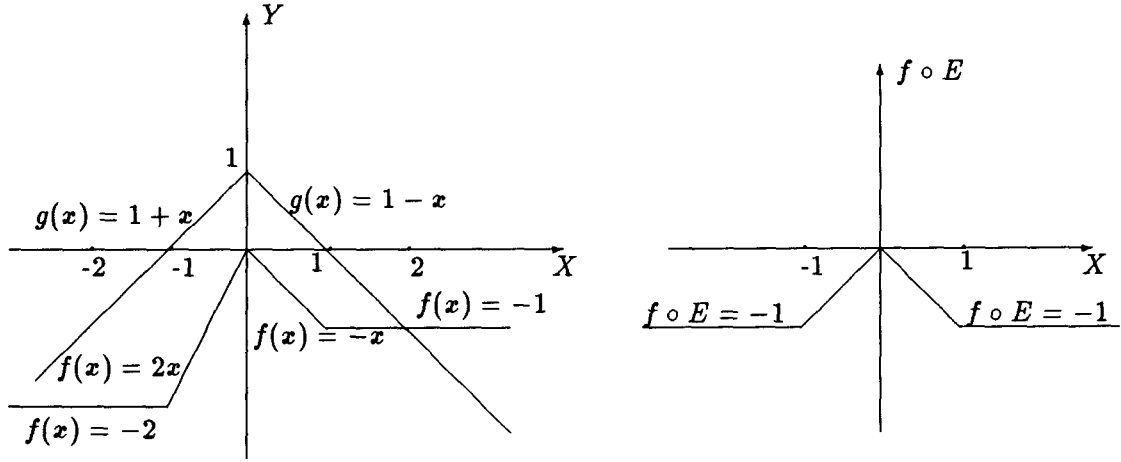


Fig.2 Example 4. A counter example for Theorem 2, i.e., Theorem 4.1 in [1]

Example 5 A counter example for Theorems 5,6, i.e., Theorems 4.4, 4.5 in [1].

Let maps $f, g, E : R \rightarrow R$ be given as

$$f(x) = \begin{cases} -1, & \text{if } |x| \geq 1; \\ -|x|, & \text{if } |x| < 1, \end{cases} \quad g(x) = -\frac{1}{2}x, \quad E(x) \equiv 0.$$

Then, see Fig. 3, we know that functions f and g all are E -convex on R , and

$$M = \{x : g(x) \leq 0\} = [0, +\infty), \quad E(M) = \{0\}; \quad (f \circ E)(x) \equiv 0, \quad (g \circ E)(x) \equiv 0, \quad \forall x \in R.$$

$$\nabla_x[(f \circ E)(x)] = \nabla_x[(g \circ E)(x)] \equiv 0, \quad \forall x \in R.$$

(i) Show that Theorem 5 is not true. From the above discusses, we know that $E(M)$ is a convex set, and f is strictly E -convex on R in the sense of Definition 2. But the solution of Problem (P_E) is not unique since each $x \in M$ is a solution of Problem (P_E) .

(ii) Show that Theorem 6 is not true. In view of the above analyses, one knows that the functions $f \circ E, g \circ E$ all are differentiable on M , and each $(x^*, y^*) \in R^2$ with $y^* \geq 0$ is a solution of Problem (4), but $E(x^*) = 0$ is not a solution of Problem (P) since its optimal solution set $\Omega = [1, +\infty)$.

Remark 2 If the strict E -convexity of f on M is defined as

$$f(\lambda E(x) + (1 - \lambda)E(y)) < \lambda f(E(x)) + (1 - \lambda)f(E(y)), \quad \forall x, y \in M, x \neq y, \lambda \in (0, 1),$$

then Theorem 5, i.e., Theorem 4.4 in [1] is true.

Example 6 A counter example for Theorem 7, i.e., Theorem 4.6 in [1].

Let maps $f, g, E : R \rightarrow R$ be given as

$$f(x) = \begin{cases} -4, & \text{if } |x| \geq 4; \\ -|x|, & \text{if } |x| < 4, \end{cases} \quad g(x) \equiv 0, \quad E(x) = \sqrt{|x|}.$$

Then, see Fig. 4, it is clear that functions f and g all are E -convex on R , and

$$M = \{x : g(x) \leq 0\} = (-\infty, +\infty) = R, \quad E(M) = E(R) = [0, +\infty).$$

But the optimal solution set $\Omega = (-\infty, -4] \cup [4, +\infty)$ is not a convex set. Furthermore Ω is not an E -convex set from Proposition 1, since $E(\Omega) = [2, +\infty) \not\subset \Omega$.

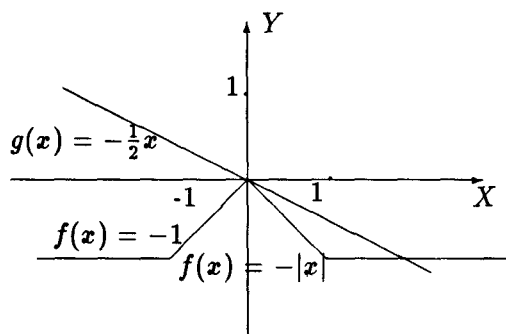


Fig.3 Example 5

A counter example for Theorems 5, 6

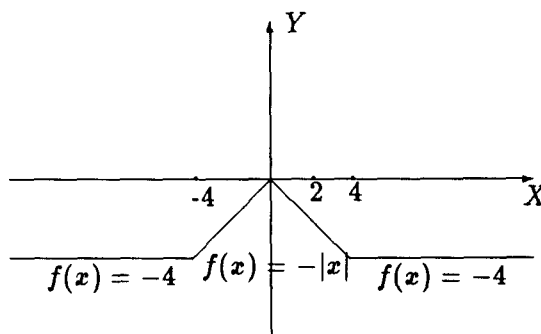


Fig.4 Example 6

A counter example for Theorem 7

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关于 E -凸函数和 E -凸规划的错误结论

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摘要: 最近 Youness 在文 [1] 建立了一类 E -凸函数和一类 E -凸规划, 并分析和给出了他们的主要性质. 本文通过 6 个反例说明文 [1] 关于 E -凸函数和 E -凸规划的大部分结论是错误的.

关键词: 凸集; 凸函数; 凸规划; 广义凸性.