

## $L(d_1, d_2, \dots, d_t)$ -Number $\lambda(C_n; d_1, d_2, \dots, d_t)$ of Cycles

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**Abstract** An  $L(d_1, d_2, \dots, d_t)$ -labeling of a graph  $G$  is a function  $f$  from its vertex set  $V(G)$  to the set  $\{0, 1, \dots, k\}$  for some positive integer  $k$  such that  $|f(x) - f(y)| \geq d_i$ , if the distance between vertices  $x$  and  $y$  in  $G$  is equal to  $i$  for  $i = 1, 2, \dots, t$ . The  $L(d_1, d_2, \dots, d_t)$ -number  $\lambda(G; d_1, d_2, \dots, d_t)$  of  $G$  is the smallest integer number  $k$  such that  $G$  has an  $L(d_1, d_2, \dots, d_t)$ -labeling with  $\max\{f(x) | x \in V(G)\} = k$ . In this paper, we obtain the exact values for  $\lambda(C_n; 2, 2, 1)$  and  $\lambda(C_n; 3, 2, 1)$ , and present lower and upper bounds for  $\lambda(C_n; 2, \dots, 2, 1, \dots, 1)$

**Keywords** cycle; labeling;  $L(d_1, d_2, \dots, d_t)$ -labeling;  $\lambda(G; d_1, d_2, \dots, d_t)$ -number.

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## 1. Introduction

The channel assignment problem is to assign a channel (nonnegative integer) to each radio transmitter so that interfering transmitters are assigned channels whose separations are not in a set of disallowable separations. Hale<sup>[1]</sup> formulated this problem into the problems of  $T$ -coloring of a graph, which has been extensively studied over the past decades<sup>[2–8]</sup>. Roberts<sup>[9]</sup> pointed that we could assign channels to some radio transmitters with different places so that close transmitters would get different channels whose difference is at least 2. Griggs and Yeh<sup>[10]</sup> first studied the problems of  $L(2, 1)$ -labeling. An  $L(2, 1)$ -labeling is a function  $f$  from its vertex set  $V(G)$  to the set  $\{0, 1, \dots, k\}$  for some integer  $k$  such that  $|f(x) - f(y)| \geq 2$  if  $d(x, y) = 1$  and  $|f(x) - f(y)| \geq 1$  if  $d(x, y) = 2$ . For positive integer numbers  $k, d_1, d_2$ , a  $k-L(d_1, d_2)$ -labeling of a graph  $G$  is a function  $f : V(G) \rightarrow \{0, 1, \dots, k\}$  such that  $|f(x) - f(y)| \geq d_i$  whenever  $x, y \in V(G)$  and  $d(x, y) = i$  ( $i = 1, 2$ ).  $L(d_1, d_2)$ -number of the graph is the smallest integer number  $k$  such that  $k-L(d_1, d_2)$ -labeling exists.

Up to now, there are a lot of results for the  $L(d_1, d_2)$ -labeling, especially, the  $L(2, 1)$ -labeling. For example, Griggs and Yeh<sup>[10]</sup> proved that the  $L(2, 1)$ -number of a tree is  $\Delta + 1$  or  $\Delta + 2$ , and that the upper bound for the  $L(2, 1)$ -number of a graph with the largest degree  $\Delta$  is at most  $\Delta^2 + 2\Delta - 3$ . Further they proposed the following conjecture is  $\Delta^2$ . In addition, they

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obtain the exact values for the  $L(2, 1)$ -number for some special graphs such as paths, cycles and wheel graphs. Chang and Kuo<sup>[11]</sup> proved that for a general graph of maximum degree  $\Delta$ , an upper bound of  $L(2, 1)$ -number is  $\Delta^2 + \Delta$ . For more background and information for the  $L(d_1, d_2)$ -numbers, the readers may refer to an excellent survey<sup>[12]</sup>.

In this survey, Yeh<sup>[12]</sup> proposed a new notion of  $L(d_1, d_2, \dots, d_t)$ -labeling of a graph. An  $L(d_1, d_2, \dots, d_t)$ -labeling of a graph  $G$  is a function  $f$  from its vertex set  $V(G)$  to the set  $\{0, 1, \dots, k\}$  for some positive integer  $k$  such that  $|f(x) - f(y)| \geq d_i$ , if the distance between vertices  $x$  and  $y$  in  $G$  is equal to  $i$  for  $i = 1, 2, \dots, t$ . The  $L(d_1, d_2, \dots, d_t)$ -number  $\lambda(G; d_1, d_2, \dots, d_t)$  of  $G$  is the smallest integer number  $k$  such that  $G$  has an  $L(d_1, d_2, \dots, d_t)$ -labeling with  $\max\{f(x) | x \in V(G)\} = k$ . Further, he proposed five problems, one of which was  $L(d_1, d_1, \dots, d_1, d_2, d_2, \dots, d_2)$ -labeling ( $d_1 > d_2 \geq 1$ ).

In this paper, we present the exact values for  $\lambda(C_n; 2, 2, 1)$ ,  $\lambda(C_n; 3, 2, 1)$  and give lower and upper bounds for  $\lambda(G; 2, 2, \dots, 2, 1, \dots, 1)$  ( $t$ -fold 2 and  $t$ -fold 1).

## 2. Preliminaries

Denote by  $C_n$  a cycle with  $n$  vertices  $v_1, v_2, \dots, v_n$ .

**Proposition 1** For a graph  $G$ , if  $\lambda(G; d_1, d_2, \dots, d_t)$  and  $\lambda(G; d_1, d_2, \dots, d_t, \delta_1, \delta_2, \dots, \delta_s)$  exist, then  $\lambda(G; d_1, d_2, \dots, d_t) \leq \lambda(G; d_1, d_2, \dots, d_t, \delta_1, \delta_2, \dots, \delta_s)$ .

**Proof** Clearly, it follows from the definition that an  $L(G; d_1, d_2, \dots, d_t, \delta_1, \delta_2, \dots, \delta_s)$ -labeling of  $G$  is also an  $L(d_1, d_2, \dots, d_t)$ -labeling. Hence the assertion holds.

**Proposition 2** For a graph  $G$ , if  $\lambda(G; d_1, d_2, \dots, d_t)$  and  $\lambda(G; \delta_1, \delta_2, \dots, \delta_t)$  exist, and  $d_i \leq \delta_i$  ( $1 \leq i \leq t$ ), then  $\lambda(G; d_1, d_2, \dots, d_t) \leq \lambda(G; \delta_1, \delta_2, \dots, \delta_t)$ .

**Proof** Since  $G$  has an  $L(\delta_1, \delta_2, \dots, \delta_t)$ -labeling,  $|f(x) - f(y)| \geq \delta_i$  for  $d(x, y) = i$  ( $1 \leq i \leq t$ ), where  $x, y \in V(G)$ . By  $d_i \leq \delta_i$  ( $1 \leq i \leq t$ ), we have  $|f(x) - f(y)| \geq d_i$ . Hence  $G$  has an  $L(d_1, d_2, \dots, d_t)$ -labeling and  $\lambda(G; d_1, d_2, \dots, d_t) \leq \lambda(G; \delta_1, \delta_2, \dots, \delta_t)$ .

**Proposition 3** Let  $G$  be a graph. If the  $L(\delta_1, \delta_2, \dots, \delta_t)$ -number exists, then there exists a vertex with labeling 0.

**Proof** Suppose the  $L(\delta_1, \delta_2, \dots, \delta_t)$ -number exists. Let the vertex  $v$  with the smallest labeling value and  $f(v) \neq 0$ . Now let  $g(u) = f(u) - f(v)$  for all  $u$  in  $G$ . Then it is easy to see that  $g$  is a function such that  $L(\delta_1, \delta_2, \dots, \delta_t)$ -number exists with  $g(v) = 0$ .

**Proposition 4** Let  $G$  be a graph with the diameter at least  $t+1$ . If the  $L(d_1, d_1, \dots, d_1, d_2, d_2, \dots, d_2)$ -labelling exists ( $t$ -fold  $d_1$ , and  $d_1 > d_2 \geq 1$ ), then  $\lambda(G; d_1, d_1, \dots, d_1, d_2, d_2, \dots, d_2) \geq td_1 + 1$ .

**Proof** Since the diameter of  $G$  is at least  $t+1$ , there exists a path with vertices  $(v_1, v_2, \dots, v_{t+1}, v_{t+2}, \dots)$ , and  $f(v_1) = 0$ . Because  $G$  has an  $L(d_1, d_1, \dots, d_1, d_2, d_2, \dots, d_2)$ -labeling, the labeling values of  $v_i$  ( $2 \leq i \leq t+1$ ) are different and  $|f(v_i) - f(v_j)| \geq d_1$  ( $1 \leq i \neq j \leq t+1$ ). Hence, among the

vertices  $v_i$  ( $2 \leq i \leq t+1$ ), there is at least one with a labeling value  $\geq td_1$ ; if the maximum of the labeling values of the vertices is  $td_1$ , then the labeling values of  $v_2, \dots, v_{t+1}$  are  $d_1, 2d_1, \dots, td_1$ . But the distance between  $v_{t+2}$  and  $v_i$  ( $2 \leq i \leq t+1$ ) is not more than  $t$  and  $f(v_{t+2}) \neq 0$ . It is impossible. So

$$\lambda(G; d_1, d_1, \dots, d_1, d_2, d_2, \dots, d_2) \geq td_1 + 1.$$

### 3. Results

**Theorem 1** For  $C_n$  ( $n \geq 3$ ), there are  $\lambda(C_n; 2, 2, 1) = \begin{cases} 4, & n = 3; \\ 8, & n = 5, 9, 13, 17; \\ 7, & n = 6, 10; \\ 6, & \text{other } n. \end{cases}$

**Proof** We first show that  $\lambda(C_n; 2, 2, 1) \geq 6$  ( $n \geq 4$ ). By Proposition 4,  $\lambda(C_n; 2, 2, 1) \geq 5$  ( $n \geq 4$ ). If  $\lambda(C_n; 2, 2, 1) = 5$  ( $n \geq 4$ ), by Proposition 3, we can set  $f(v_1) = 0$ , thus  $f(v_2) \geq 2$ ; if  $2 \leq f(v_2) \leq 4$ , with  $f(v_3) \geq 2$  and  $|f(v_2) - f(v_3)| \geq 2$ , then  $f(v_3) \in \{2, 4, 5\}$ . Hence if  $n = 4$ , then there are no labeling values for  $v_4$ ; if  $n = 5$ , then there are no labeling values for  $v_4$ ; if  $n \geq 6$  with  $|f(v_2) - f(v_4)| \geq 2, |f(v_3) - f(v_4)| \geq 2$ , then there is only one labeling:  $f(v_2) = 3, f(v_3) = 5, f(v_4) = 1$ , but, there is no labeling value for  $v_5$ . If  $f(v_2) = 5$ , then  $f(v_3) \in \{2, 3\}$ , but  $f(v_4) \geq 1$ . So we have only one labeling, that is,  $f(v_2) = 5, f(v_3) = 3, f(v_4) = 1$ . In this case, there is no labeling value for  $v_5$  either. Therefore  $\lambda(C_n; 2, 2, 1) \geq 6$  ( $n \geq 4$ ).

Now we can obtain the results by constructing labeling. If  $n = 3$ , set  $f : v_1 v_2 v_3 \rightarrow 024$ ; if  $n = 5$ , set  $f : v_1 v_2 \cdots v_5 \rightarrow 02468$ ; if  $n = 9$ , set  $f : v_1 v_2 \cdots v_9 \rightarrow 024681357$ ; if  $n = 13, 17$ , the first 9 vertices are valued as  $n = 9$ , the left vertices are valued as 0246, 02460246; if  $n = 6$ , set  $f : v_1 v_2 \cdots v_6 \rightarrow 037146$ ; if  $n = 10$ , the first 6 vertices are valued as  $n = 6$ , the left vertices are valued as 0246.

Finally, we show that  $\lambda(C_n; 2, 2, 1) = 6$  for the remaining case. We now construct the following labeling.

If  $n \equiv 0 \pmod{4}$ , set  $f : v_1 \cdots v_4 \rightarrow 0246, f(v_i) = f(v_{i+4})$ .

If  $n \equiv 1 \pmod{4}$  and  $n > 17$ , set  $f : v_1 \cdots v_7 \rightarrow 0246135, f(v_i) = f(v_{i+7})$ , where  $i = 1, 2, \dots, 14$ ; for the left vertices, the labeling is as  $n \equiv 0 \pmod{4}$ .

If  $n \equiv 2 \pmod{4}$  and  $n > 10$ , set  $f : v_1 \cdots v_7 \rightarrow 0246135, f(v_i) = f(v_{i+7})$ , where  $i = 1, 2, \dots, 7$ ; for the left vertices, the labeling is as  $n \equiv 0 \pmod{4}$ .

If  $n \equiv 3 \pmod{4}$ , set  $f : v_1 \cdots v_7 \rightarrow 0246135$ , for the left vertices, the labeling is as  $n \equiv 0 \pmod{4}$ .

By simple calculations, it is easy to see that labeling of the above is  $L(2, 2, 1)$ -labeling of cycle  $C_n$ .  $\square$

**Theorem 2** For  $C_n$  ( $n \geq 3$ ), there are (1)  $\lambda(C_n; 3, 2, 1) = \begin{cases} 6, & n = 3; \\ 9, & n = 7. \end{cases}$

(2)  $\lambda(C_n; 3, 2, 1) = \begin{cases} 8, & n > 3 \text{ (} n \neq 7 \text{), and is odd;} \\ 7, & n \geq 4, \text{ and is even.} \end{cases}$

**Proof** By Proposition 1 and some calculations, it is easy to see that  $\lambda(C_n; 3, 2, 1) = 6, 7, 8, 7, 9$ , corresponding to  $n = 3, 4, \dots, 7$ , respectively.

Now we assume that  $n > 7$ . If  $\lambda(C_n; 3, 2, 1) = 6$ ,  $f(v_1) = 0$ ; if  $f(v_2) = 3$ , then  $f(v_3) = 6$ ,  $f(v_4) = 1$ ,  $f(v_5) = 4$ , but there is also no labeling value for  $v_6$ ; if  $f(v_2) = 4$ , there is no labeling value for  $v_3$ ; if  $f(v_2) = 5$  or  $6$ , then  $f(v_3) = 2$  or  $3$ , there is no labeling value for  $v_4$ . Hence  $\lambda(C_n; 3, 2, 1) \geq 7$ .

If  $n \geq 8$ , and is even, set  $f$  :

$$v_1 \cdots v_4 \rightarrow 0725, f(v_i) = f(v_{i+4}) \ (i \geq 1), \text{ if } n \equiv 0 \pmod{4};$$

$$v_1 \cdots v_6 \rightarrow 036147, f(v_i) = f(v_{i+6}) \ (i \geq 1), \text{ if } n \equiv 0 \pmod{6};$$

$$v_1 \cdots v_6 \rightarrow 036147, f(v_i) = f(v_{i+6}) \ (1 \leq i \leq n-8), v_{n-1}v_n \rightarrow 25, \text{ if } n \equiv 2 \pmod{6};$$

$v_1 \cdots v_6 \rightarrow 036147, f(v_i) = f(v_{i+6}) \ (1 \leq i \leq n-10), v_{n-3} \cdots v_n \rightarrow 0527$ , if  $n \equiv 4 \pmod{6}$ , and  $n \neq 4k$ . Hence  $\lambda(C_n; 3, 2, 1) = 7$ .

If  $n \geq 9$ , and is odd, we show that  $\lambda(C_n; 3, 2, 1) = 8$ . In fact, if  $\lambda(C_n; 3, 2, 1) = 7$ , for the labelings that can be recirculated on  $C_n$  are: 0725; 036147; 03614725, the number in each set is even, and each labeling can be removed. So if we label the vertices of  $C_n$  by use of these sets, we cannot label the remaining odd vertices of  $C_n$  by use of the numbers in  $\{0, 1, \dots, 7\}$ . Thus, we obtain  $\lambda(C_n; 3, 2, 1) > 7$  when  $n \geq 9$ , and is odd.

If  $n \equiv 3 \pmod{4}$ , set  $f$ :  $v_1 \cdots v_7 \rightarrow 0741836, v_8 \cdots v_{11} \rightarrow 0825, f(v_i) = f(v_{i+4}) \ (i \geq 8)$ ; If  $n \equiv 1 \pmod{4}$ , set  $f$ :  $v_1 \cdots v_5 \rightarrow 04826, v_8 \cdots v_{11} \rightarrow 0826, f(v_i) = f(v_{i+4}) \ (i \geq 6)$ . Therefore  $\lambda(C_n; 3, 2, 1) = 8$ .  $\square$

**Theorem 3** For  $C_n (n \geq 3)$ , there are  $\lambda(C_n; 2, \dots, 2, 1, \dots, 1) \leq 4t$  ( $t$ -fold 2 and 1).

**Proof** From the proofs of Theorems 1 and 2, we see that the key step in labeling a cycle is how to construct the labeling of  $C_n \ (3 \leq n \leq 4t)$ . We will do this.

If  $3 \leq n \leq 2t+1$ , set  $f(V) \rightarrow 024 \cdots (2n-2)$ .

In case of  $2t+2 \leq n \leq 4t$ :

(1) If  $n$  is odd, set

$$f(V) \rightarrow 0(4t)(4t-2) \cdots (4t-2[\frac{n}{2}] + 2)1(4t-1)(4t-3) \cdots (4t-2[\frac{n}{2}] + 3);$$

(2) If  $n$  is even, set

$$f(V) \rightarrow 0(4t)(4t-2) \cdots (4t-2[\frac{n}{2}] + 4)1(4t-1)(4t-3) \cdots (4t-2[\frac{n}{2}] + 3).$$

If  $n = 4t+1$ , set

$$f(V) \rightarrow 0(4t)(4t-2) \cdots (2t+2)(2t-1)(2t-3) \cdots 31(4t-1)(4t-3) \cdots 9(2t+1)24 \cdots (2t).$$

If  $n > 4t+1$ , and  $n \neq k(2t+1)$ ,  $k = 2, 3, \dots$ , we separate the vertices of  $C_n$  into two parts. The number of the vertices in one part is a multiple of  $2t+1$ , which are circularly labeled by  $0(4t)(4t-2) \cdots 2$ . The number of the vertices in other part is between  $2t+2$  and  $4t+1$ , which are labeled in the same way as the above case of  $2t+2 \leq n \leq 4t+1$ . So  $L(2, \dots, 2, 1, \dots, 1)$ -labeling exists.  $\square$

**Theorem 4** For  $C_n(n \geq 3)$ , there are  $\lambda(C_n; 2, 1, \dots, 1) = \begin{cases} 4, & n = 3, 4; \\ 2t + 2, & \text{other } n(t\text{-fold } 1). \end{cases}$

**Proof** If  $n \leq 2t + 3$ , the diameter of the cycle is  $\lceil \frac{n}{2} \rceil$ , so the labeling values of vertices are different from each other and  $\lambda(C_n; 2, 1, \dots, 1)$  must be more than  $n - 1$ . Next we will show that except for  $n = 3, 4$ ,  $\lambda(C_n; 2, 1, \dots, 1) = n - 1 \leq 2t + 2$ .

(1)  $n = 3$ , set  $f(V) \rightarrow 024$ ;  $n = 4$ , set  $f(V) \rightarrow 0314$ .

(2) If  $4 \leq n \leq 2t + 3$ , we will do the labeling by the following rule: if  $n$  is even, set  $f(V) \rightarrow 024 \cdots (n - 2)13 \cdots (n - 1)$ ; if  $n$  is odd, set  $f(V) \rightarrow 024 \cdots (n - 1)13 \cdots (n - 2)$ . So, for  $C_n$ , it is obvious that  $L(2, 1, \dots, 1)$ -labeling exists.

If  $n = k(2t + 4)$   $k = 1, 2, \dots$ , set

$$f(V) \rightarrow 024 \cdots (2t + 2)024 \cdots (2t + 2)024 \cdots (2t + 2).$$

If  $n > 2t + 3$  ( $n \neq k(2t + 4), k = 1, 2, \dots$ ), we separate the vertices of  $C_n$  into two parts, the number of the vertices in one part is a multiple of  $t + 2$ , and the number of the vertices in the other part is between  $t + 3$ , and  $2t + 3$ . For the first part, if  $t + 1$  is even, the vertices are circularly labeled by  $024 \cdots (t + 1)13 \cdots t$ ; if  $t + 1$  is odd, the vertices are circularly labeled by  $024 \cdots t13 \cdots (t + 1)$ . For the other part, the vertices are labeled in the same way as the above. It can be proven that for any  $n$  and any cycle, by the above labeling  $L(2, 1, \dots, 1)$ -labeling exists.  $\square$

**Corollary** For  $C_n(n \geq 3)$ , there are  $2t + 2 \leq \lambda(C_n; 2, \dots, 2, 1) \leq 4t$  ( $t$ -fold 2).

**Proof** By Propositions 1 and 2, the assertion holds.

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