# The Factor Spectrum and Derived Sequence 

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Dedicated to the Memory of Professor L. C. HSU on the Occasion of His 100th Birthday


#### Abstract

Given a sequence $\rho$ over a finite alphabet $\mathcal{A}$, an important topic in combinatorics on words is to find out all factors $\omega$ of $\rho$ and positive integers $p$ such that $\omega_{p}$ (the $p$-th occurrence of $\omega$ ) fulfills property $\mathcal{P}$. This problem is equivalent to determining a notion called the factor spectrum. Determining the factor spectrum is a difficult problem. To this aim, we introduce several notions, such as: kernel word, envelope word, return word and derived sequence of each factor $\omega$. Using the factor spectrum and derived sequence, we can solve some enumerations of factors, such as the numbers of palindromes, fractional powers, etc. We will show some results for several sequences, such as the Fibonacci sequence, the Tribonacci sequence, the Period-doubling sequence, etc. And we think that these notions and methods are suitable for all recurrent sequences.


Keywords kernel word; envelope word; return word; derived sequence; the factor spectrum
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## 1. Introduction

A sequence $\rho$ is said to be recurrent if every factor occurs infinitely often [1]. We arrange all occurrences of factor $\omega$ in $\rho$ as the sequence $\left\{\omega_{p}\right\}_{p \geq 1}$, where $\omega_{p}$ denotes the $p$-th occurrence of $\omega$. Let $\mathcal{P}$ be a property about factor, and if the $p$-th occurrence of $\omega$ has property $\mathcal{P}$, we denote $\omega_{p} \in \mathcal{P}$. We denote $\omega \in \mathcal{P}$, if there exists an integer $p \geq 1$ such that $\omega_{p} \in \mathcal{P}$. When $\omega \in \mathcal{P}$, maybe not all $p \in \mathbb{N}$ such that $\omega_{p} \in \mathcal{P}$. An important topic in combinatorics on words is to find out all $\omega$ and $p$ such that $\omega_{p}$ fulfills property $\mathcal{P}$. More precisely, let $\rho$ be a sequence and $\mathcal{P}$ be a property, we introduce the following notion factor spectrum of a sequence $\rho$ and a property $\mathcal{P}$ that

$$
\begin{equation*}
\operatorname{Spt}(\rho, \mathcal{P}):=\left\{(\omega, p) \mid \omega \prec \rho, p \geq 1, \omega_{p} \in \mathcal{P}\right\} \subset\left(\Omega_{\rho}, \mathbb{N}\right) \tag{1.1}
\end{equation*}
$$

where $\Omega_{\rho}$ is the set of all factors in $\rho$. When we write $(\omega, p) \in \operatorname{Spt}(\rho, \mathcal{P})$, we regard $\omega$ and $p$ as two variable function over $\left(\Omega_{\rho}, \mathbb{N}\right)$. By the definition in Eq. (1.1), the above problem is equivalent to determining the factor spectra.

[^0]Determining the factor spectra is a difficult problem. For this aim, we use the return word and derived sequence of each factor $\omega$, which are introduced in Durand [2], see Section 2 for details. From them, we may know the local structure near $\omega_{p}$ for each $p$. In Huang-Wen [3, 4], we first study the structure of derived sequences of the Fibonacci sequence and the Tribonacci sequence, then we determine the factor spectra for some combinatorial properties.

## 2. The return word and derived sequence

Let $\omega=x_{1} x_{2} \cdots x_{N}$ be a finite word. We denote by $|\omega|$ the number of letters in $\omega$, called the length of $\omega$. For $1 \leq i \leq j \leq N$, we define $\omega[i, j]=x_{i} x_{i+1} \cdots x_{j-1} x_{j}$, by convention, $\omega[i]=\omega[i, i]=x_{i}$ and $\omega[i, i-1]=\varepsilon$ (empty word). We call $\omega[i, j]$ a factor of $\omega$, denoted by $\omega[i, j] \prec \omega$. The position of factor $\omega[i, j]$ in $\omega$ is defined by $i$. We call a non-empty word $\nu$ a prefix (resp., suffix) of a word $\omega$ if there exists a word $u$ such that $\omega=\nu u$ (resp., $\omega=u \nu$ ), denoted by $\nu \triangleleft \omega$ (resp., $\nu \triangleright \omega$ ). In this case, we write $\nu^{-1} \omega=u$ (resp., $\omega \nu^{-1}=u$ ), where $\nu^{-1}$ is the inverse word of $\nu$ such that $\nu \nu^{-1}=\nu^{-1} \nu=\varepsilon$. A palindrome $\omega$ is a finite word that reads the same backwards as forwards, i.e., $\overleftarrow{\omega}:=x_{N} \cdots x_{2} x_{1}=\omega$.

The definitions of both the return word and derived sequence are from Durand [2]. Let us recall them as below. Notice that Durand gave these notions for all prefixes of any recurrent sequence, and we extend these notions to all non-empty factors. Recall that $\omega_{p}$ denotes the $p$-th occurrence of $\omega$. Denote by $\operatorname{occ}(\omega, p)$ the position of $\omega_{p}$. For $p, q \geq 1$, $\omega_{p} \prec W_{q}$ means $\omega \prec W$ and $\operatorname{occ}(W, q) \leq \operatorname{occ}(\omega, p)<\operatorname{occ}(\omega, p)+|\omega|-1 \leq \operatorname{occ}(W, q)+|W|-1$.

We call $\rho[i, j-1]$ the $p$-th return word over $\omega$ where $i$ and $j$ are the positions of the $p$-th and $(p+1)$-th occurrences of $\omega$ in $\rho$, denoted by $R_{\rho, p}(\omega)$. Denote by $\mathcal{H}_{\rho, \omega}$ the set of return words over factor $\omega \prec \rho$. Then the sequence $\rho$ can be written in a unique way as a concatenation $\rho=\rho[1, h-1] R_{\rho, 1}(\omega) R_{\rho, 2}(\omega) \cdots$ where $R_{\rho, p}(\omega) \in \mathcal{H}_{\rho, \omega}$ and $\rho[1, h-1]$ is the prefix of $\rho$ occurring before the first occurrence of $\omega$. Let us give to $\mathcal{H}_{\rho, \omega}$ the linear order defined by the rank of the first occurrence in $\rho$. This defines a one to one and onto map $\Lambda_{\rho, \omega}: \mathcal{H}_{\rho, \omega} \rightarrow\left\{1, \ldots, \operatorname{Card}\left(\mathcal{H}_{\rho, \omega}\right)\right\}=$ $\mathcal{N}_{\rho, \omega} \subset\{\alpha, \beta, \ldots\}$, and the sequence $\mathcal{D}_{\omega}(\rho):=\Lambda_{\rho, \omega}\left(R_{\rho, 1}(\omega)\right) \Lambda_{\rho, \omega}\left(R_{\rho, 2}(\omega)\right) \Lambda_{\rho, \omega}\left(R_{\rho, 3}(\omega)\right) \cdots$. This sequence of alphabet $\mathcal{N}_{\rho, \omega}$ is called a derived sequence of $\rho$. Notice that we omit the prefix $\rho[1, h-1]$. Moreover, we denote the reciprocal map of $\Lambda_{\rho, \omega}$ by $\Theta_{\rho, \omega}: \mathcal{N}_{\rho, \omega} \rightarrow \mathcal{H}_{\rho, \omega}$.

The main result of Durand [2] is: a sequence $\rho$ is substitutive primitive if and only if the number of its different derived sequences is finite. The other property is: for any $\omega \prec \rho$ and $v \prec \mathcal{D}_{\omega}(\rho)$, there exists a factor $\mu \prec \rho$ such that derived sequence $\mathcal{D}_{v}\left(\mathcal{D}_{\omega}(\rho)\right)=\mathcal{D}_{\mu}(\rho)$ (see [2, Proposition $6(5)]$ ). By the two properties, if $\rho$ is substitutive primitive, then for any $\omega \prec \rho$, derived sequence $\mathcal{D}_{\omega}(\rho)$ is still substitutive primitive.

### 2.1. Derived sequences of the Fibonacci and Tribonacci sequences

The Fibonacci morphism $\sigma_{1}$ over alphabet $\{a, b\}$ is a substitution defined by $\sigma_{1}(a)=a b$ and $\sigma_{1}(b)=a$. The Fibonacci sequence $\mathbb{F}$ is defined to be the fixed point beginning with the letter $a$ of the Fibonacci morphism. It is a recurrent infinite word [1]. Define $F_{m}=\sigma_{1}^{m}(a)$ for $m \geq 0$,
by convention, $F_{-1}=b$ and $F_{-2}=\varepsilon$. The $m$-th Fibonacci number $f_{m}$ is equal to the length of $F_{m}$. As a classical example of sequences over the binary alphabet, the Fibonacci sequence $\mathbb{F}$ has many remarkable properties. We refer to Lothaire [5, 6], Allouche-Shallit [1] and Berstel [7, 8]. The Tribonacci sequence $\mathbb{T}$ is a natural generalization of the Fibonacci sequence, which is the fixed point beginning with the letter $a$ of morphism $\sigma_{2}(a)=a b, \sigma_{2}(b)=a c, \sigma_{2}(c)=a$ defined over the alphabet $\{a, b, c\}$. Define $T_{m}=\sigma_{2}^{m}(a)$ for $m \geq 0$.

The $m$-th Tribonacci number $t_{m}$ is equal to the length of $T_{m}$.

### 2.1.1. Kernel word

The main tool of Huang-Wen [3,4] is kernel word, which is a set of factors in the sequence $\rho$. For a factor $\omega$, we first introduce the kernel of a factor $\omega$ as follows, denoted by $\operatorname{Ker}(\omega)$.

$$
\begin{align*}
\operatorname{Ker}(\omega):=\min _{|\cdot|}\{W \mid & \text { the difference } \operatorname{occ}(W, p)-\operatorname{occ}(\omega, p) \\
& \text { is independent of } p \text { for } p \geq 1, W \prec \omega\} . \tag{2.1}
\end{align*}
$$

We can prove that the kernel of any factor $\omega$ is unique when $\rho$ is the Fibonacci or Tribonacci sequence. If $\operatorname{Ker}(\omega)$ is unique for all $\omega$ in $\rho$, we define the set of kernel words by that $\mathcal{K}:=$ $\{\operatorname{Ker}(\omega) \mid \omega \prec \rho\}$. More precisely, we give the definitions of kernel words in the Fibonacci and Tribonacci sequences as below. In fact, the kernel word in the Fibonacci sequence is just the singular words defined in Wen-Wen [9]. But unfortunately, the singular words of the Tribonacci sequence defined in Tan-Wen [10] is not a set satisfying Eq. (2.1), see Example 2.2(b).

Remark 2.1 The two notions are suitable not only for the Fibonacci and Tribonacci sequences, but also for some other sequences. Notice that the notions are not suitable for the Perioddoubling sequence $\mathbb{D}$, for details see Example 2.2(a). For this kind of sequences, we need other tools.

Example 2.2 (a) Take $\omega=b a a a b a b a b a a a b \prec \mathbb{D}, U=a a b a b$ and $V=b a b a a$. We can check that $|U|=|V|, \operatorname{occ}(U, p)-\operatorname{occ}(\omega, p) \equiv 2$ and $\operatorname{occ}(V, p)-\operatorname{occ}(\omega, p) \equiv 6$ for all $p \geq 1$. But for any factor $W$ in $\mathbb{D}$ of length $1 \leq n \leq 4$, it dose not satisfy that $\operatorname{occ}(W, p)-\operatorname{occ}(\omega, p)$ is independent of $p$. This means both $U$ and $V$ satisfy Eq. (2.1).
(b) Tan-Wen [10] defined two kinds of singular words: $\Omega_{m}^{1}$ and $\Omega_{m}^{2}$ for $m \geq 1$.

$$
\left\{\begin{array}{l}
\Omega_{m}^{1}=\left\{\text { factor of length } f_{m} \text { of the word } \alpha^{-1} \overleftarrow{E_{m}} D_{m-1} E_{m} \alpha^{-1}\right\} \\
\Omega_{m}^{2}=\left\{\text { factor of length } f_{m} \text { of the word } \beta^{-1} E_{m+1} D_{m-2} \overleftarrow{E_{m+1}} \beta^{-1}\right\}
\end{array}\right.
$$

where $D_{m}=T_{m-1} T_{m-2} \cdots T_{2} T_{1} T_{0}, E_{m}=D_{m-1}^{-1} T_{m}, \alpha$ is the last letter of $E_{m}$ and $\beta$ is the first letter of $E_{m+1}$. By convention, $D_{0}=\varepsilon$. For instance, $\Omega_{1}^{1}=\{a a\}, \Omega_{1}^{2}=\{a c, c a\}, \Omega_{2}^{1}=$ $\{a b a b, b a b a\}$ and $\Omega_{2}^{2}=\{a b a a, b a a b, a a b a\}$. We can check that the set of all singular words in $\mathbb{T}$ is not satisfying Eq. (2.1).

Definition 2.3 (Kernel words of the Fibonacci sequence) (1) Let $\left\{K_{m}\right\}_{m \geq-1}$ be the sequence
of factors with

$$
K_{-1}=a \text { and } K_{m}=\delta_{m+1} F_{m}\left[1, f_{m}-1\right] \text { for } m \geq 0
$$

where $\delta_{m} \triangleright F_{m}$. We call $K_{m}$ the $m$-th kernel word. More precisely, $\delta_{m}=a$ if $m \equiv 0(\bmod 2)$, $\delta_{m}=b$ otherwise. Obviously, $\left|K_{m}\right|=f_{m}$ for $m \geq-1$.
(2) Define $\mathcal{K}:=\left\{K_{m} \mid m \geq-1\right\}$, which is called the kernel set of the Fibonacci sequence.
(3) We define the order of kernel words that $K_{m} \sqsubset K_{m+1}$ for all $m \geq-1$. For any factor $\omega \prec \mathbb{F}$, we define

$$
\operatorname{Ker}(\omega)=\max _{\sqsubset}\left\{K_{m} \mid K_{m} \prec \omega, m \geq-1\right\}
$$

which is called the kernel of factor $\omega$.
(4) We define the order $K_{m, p} \sqsubset K_{m+1, q}$ for all $p, q \geq 1$. For any factor $\omega_{p}$ where $(\omega, p) \in$ $\left(\Omega_{\mathbb{F}}, \mathbb{N}\right)$, we define

$$
\operatorname{Ker}\left(\omega_{p}\right)=\max _{\sqsubset}\left\{K_{m, q} \mid K_{m, q} \prec \omega_{p}, m \geq-1, q \geq 1\right\}
$$

which is called the kernel of factor $\omega_{p}$.
Here are the first few values of $F_{m}, f_{m}$ and $K_{m}$ :

- $\left\{F_{m}\right\}_{m \geq-1}=b, a, a b, a b a, a b a a b, a b a a b a b a, a b a a b a b a a b a a b, a b a a b a b a a b a a b a b a a b a b a, \ldots$..
- $\left\{f_{m}\right\}_{m \geq-1}=1,1,2,3,5,8,13,21,34,55,89,144,233,377,610,987,1597,2584,4181, \ldots$.
- $\left\{K_{m}\right\}_{m \geq-1}=a, b, a a, b a b, a a b a a, b a b a a b a b, a a b a a b a b a a b a a$, babaababaabaababaabab, $\ldots$.

We have proved that all kernel words of $\mathbb{F}$ are palindrome [3].
Definition 2.4 (Kernel words of the Tribonacci sequence) (1) Let $\left\{k_{m}\right\}_{m \geq 1}$ be the sequence of positive integers with

$$
k_{1}=k_{2}=k_{3}=1 \text { and } k_{m}=k_{m-1}+k_{m-2}+k_{m-3}-1 \text { for } m \geq 4
$$

The number $k_{m}$ is called the $m$-th kernel number.
(2) Let $\left\{K_{m}\right\}_{m \geq 1}$ be the sequence of factors with

$$
K_{1}=a, K_{2}=b, K_{3}=c \text { and } K_{m}=\delta_{m} T_{m-3}\left[1, k_{m}-1\right] \text { for } m \geq 4
$$

where $\delta_{m} \triangleright T_{m}$. We call $K_{m}$ the $m$-th kernel word. More precisely, $\delta_{m}=a$ if $m \equiv 0(\bmod 3)$, $\delta_{m}=b$ if $m \equiv 1(\bmod 3), \delta_{m}=c$ otherwise. Obviously, $\left|K_{m}\right|=k_{m}$ for $m \geq 1$.
(3) Define $\mathcal{K}:=\left\{K_{m} \mid m \geq 1\right\}$, which is called the kernel set of the Tribonacci sequence.
(4) We define the order of kernel words that $K_{m} \sqsubset K_{m+1}$ for all $m \geq 1$. For any factor $\omega \prec \mathbb{T}$, we define $\operatorname{Ker}(\omega)=\max _{\sqsubset}\left\{K_{m} \mid K_{m} \prec \omega, m \geq 1\right\}$, which is called the kernel of factor $\omega$.
(5) We define the order $K_{m, p} \sqsubset K_{m+1, q}$ for $m, p, q \geq 1$. For any factor $\omega_{p}$ where $(\omega, p) \in$ $\left(\Omega_{\mathbb{T}}, \mathbb{N}\right)$, we define $\operatorname{Ker}\left(\omega_{p}\right)=\max _{\sqsubset}\left\{K_{m, q} \mid K_{m, q} \prec \omega_{p}, m, q \geq 1\right\}$, which is called the kernel of factor $\omega_{p}$.

Here are the first few values of $T_{m}, t_{m}, K_{m}$ and $k_{m}$ :

- $\left\{T_{m}\right\}_{m \geq 0}=a, a b, a b a c, a b a c a b a, a b a c a b a a b a c a b, a b a c a b a a b a c a b a b a c a b a a b a c, \ldots$.
- $\left\{t_{m}\right\}_{m \geq 0}=1,2,4,7,13,24,44,81,149,274,504,927,1705,3136, \ldots$.
- $\left\{K_{m}\right\}_{m \geq 1}=a, b, c, a a, b a b, c a b a c, a a b a c a b a a, b a b a c a b a a b a c a b a b, \ldots$.
- $\left\{k_{m}\right\}_{m \geq 1}=1,1,1,2,3,5,9,16,29,53,97,178,327,601, \ldots$.

As the case of the Fibonacci sequence, all kernel words of $\mathbb{T}$ are also palindrome [4].
Remark 2.5 Notice that, the definitions of $\operatorname{Ker}(\omega)$ in Eq. (2.1) and Definitions 2.3 and 2.4 look different. In fact, they are equivalent. More precisely, Eq. (2.1) is more essential, and Definitions 2.3 and 2.4 are more applicable and more convenient.

### 2.1.2. Derived sequences

Huang-Wen proved the properties about "derived sequence" below, see [3, Theorem 2.11] and [4, Theorem 5.1], respectively.

Theorem 2.6 (1) For any factor $\omega \prec \mathbb{F}$, derived sequence $\mathcal{D}_{\omega}(\mathbb{F})$ is still $\mathbb{F}$ itself.
(2) For any factor $\omega \prec \mathbb{T}$, derived sequence $\mathcal{D}_{\omega}(\mathbb{T})$ is still $\mathbb{T}$ itself.

Using Theorem 2.6, we determined the numbers of palindromes, squares and cubes occurring in $\mathbb{F}[i, j]$ and $\mathbb{T}[i, j]$ for $1 \leq i \leq j$, see Huang-Wen [11, 12] for instance. These topics are of great importance in computer science.

Example 2.7 (Derived sequence $\mathcal{D}_{\omega}(\mathbb{F})$ for $\omega=b a a b$ )

$$
\mathbb{F}=a \underbrace{b a a b a}_{\alpha} \underbrace{b a a}_{\beta} \underbrace{b a a b a}_{\alpha} \underbrace{b a a b a}_{\alpha} \underbrace{b a a}_{\beta} \underbrace{b a a b a}_{\alpha} \underbrace{b a a}_{\beta} \underbrace{b a a b a}_{\alpha} \underbrace{b a a b a}_{\alpha} \underbrace{b a a}_{\beta} \underbrace{b a a b a}_{\alpha} \underbrace{b a a b a}_{\alpha} \underbrace{b a a}_{\beta} \cdots,
$$

where

$$
\alpha=\Lambda_{\mathbb{F}, \omega}\left(R_{\mathbb{F}, 1}(\omega)\right)=\Lambda_{\mathbb{F}, \omega}(b a a b a) \text { and } \beta=\Lambda_{\mathbb{F}, \omega}\left(R_{\mathbb{F}, 2}(\omega)\right)=\Lambda_{\mathbb{F}, \omega}(b a a)
$$

It is easy to see that the derived sequence $\mathcal{D}_{\omega}(\mathbb{F})=\alpha \beta \alpha \alpha \beta \alpha \beta \alpha \alpha \beta \alpha \alpha \beta \cdots$ is a Fibonacci sequence over $\{\alpha, \beta\}$.

### 2.1.3. Three steps to prove Theorem 2.6

How to prove Theorem 2.6 by the kernel word? We need three steps:
Step 1. Determine derived sequence $\mathcal{D}_{\omega}(\rho)$ for any kernel words $\omega \prec \mathcal{K}$.
Step 2. For all $\omega \prec \rho$, prove that the difference $\operatorname{occ}(\operatorname{Ker}(\omega), p)-\operatorname{occ}(\omega, p)$ is independent of p.

Step 3. Prove derived sequence $\mathcal{D}_{\omega}(\rho)$ is exactly $\mathcal{D}_{\operatorname{Ker}(\omega)}(\rho)$.
For details, see Huang-Wen [3, 4] for instance.


Figure 1 The relation between $\omega_{p}$ and $\operatorname{Ker}(\omega)_{p}$ for $\omega=b a a b$ and $1 \leq p \leq 5$.

### 2.2. Derived sequences of the Period-doubling sequence

Huang-Wen [3,4] determine the structure of derived sequences of the Fibonacci sequence and the Tribonacci sequence, respectively. But for the Period-doubling sequence $\mathbb{D}$, there are two differences. Firstly, we find that for different factors, they may have different derived sequences. But fortunately, we can divide set $\Omega_{\mathbb{D}}$ into two types, and each type corresponds to one derived sequence. Secondly, the main tool of $[3,4]$ is kernel word. But the kernel word technique is not valid for the Period-doubling sequence, thus we introduce the envelope words and corresponding techniques. In this way, we determine the factor spectra for some combinatorics properties.

Let $\mathcal{A}=\{a, b\}$ be a binary alphabet. The Period-doubling sequence $\mathbb{D}$ is the fixed point beginning with the letter $a$ of substitution $\sigma_{3}(a)=a b$ and $\sigma_{3}(b)=a a$. It is also the first difference of the Thue-Morse sequence, where we use an equivalent substitution $\hat{\sigma}(1)=10, \hat{\sigma}(0)=11$, and the definition of the difference of an integer sequence is natural. It has been heavily studied in mathematics and computer science. Damanik [13] determined the numbers of palindromes, squares and cubes of length $n$ occurring in $\mathbb{D}$. Allouche-Peyrière-Wen-Wen [14] proved that all the Hankel determinants of $\mathbb{D}$ are odd integers [15-20, 22-28]. We denote $A_{m}=\sigma_{3}^{m}(a)$ and $B_{m}=\sigma_{3}^{m}(b)$ for $m \geq 0$. Then $\left|A_{m}\right|=\left|B_{m}\right|=2^{m}$. Let $\delta_{m} \in\{a, b\}$ be the last letter of $A_{m}$. Obviously, $\delta_{m}=a$ if and only if $m$ is even; and $\delta_{m+1}$ is the last letter of $B_{m}$.

Recall that for any factor $\omega$ of $\mathbb{F}$, derived sequence $\mathcal{D}_{\omega}(\mathbb{F})$ is still $\mathbb{F}$ itself. But HuangWen [29] proved that for different factors in $\mathbb{D}$, there will be different derived sequences. In [29], we determined two types of derived sequences, see Theorem 2.9.

### 2.2.1. Envelope word

The main tool of Huang-Wen [29] is envelope word, which is a set of factors in sequence. We first introduce the envelope of a factor $\omega$ for the factor in sequence $\rho$, denoted by $\operatorname{Env}(\omega)$.

$$
\begin{align*}
\operatorname{Env}(\omega):=\max _{|\cdot|}\{W \mid & \text { the difference } \operatorname{occ}(\omega, p)-\operatorname{occ}(W, p) \\
& \text { is independent of } p \text { for } p \geq 1, \omega \prec W\} . \tag{2.2}
\end{align*}
$$

We can prove that the envelope of any factor $\omega$ is unique. In fact, if $U$ and $V$ are two words satisfying Eq. (2.2) and $\operatorname{occ}(U, 1)<\operatorname{occ}(V, 1)$, then $W=\rho[\operatorname{occ}(U, 1), \operatorname{occ}(V, 1)+|V|-1]$ also satisfies Eq. (2.2). Since $|W|>|U|=|V|$, a contradiction. We define the set of envelope words by that $\mathcal{E}:=\{\operatorname{Env}(\omega) \mid \omega \prec \rho\}$.

More precisely, we give the definitions of envelope words in the Period-doubling sequence as below. Notice that, the definitions of $\operatorname{Env}(\omega)$ in Eq. (2.2) and Definition 2.8 look different. In fact, they are equivalent. More precisely, Eq. (2.2) is more essential, and Definition 2.8 is more applicable and more convenient.

We think notions of envelop words will be suitable for all recurrent sequence.
Definition 2.8 (Envelope word of the Period-doubling sequence) (1) Let $\left\{E_{m}^{i} \mid i=1,2, m \geq 1\right\}$
be a set of factors with

$$
E_{m}^{1}=A_{m} \delta_{m}^{-1} \text { and } E_{m}^{2}=B_{m} B_{m-1} \delta_{m}^{-1}
$$

We call $E_{m}^{i}$ the m-th envelope word of type $i$. Moreover $\left|E_{m}^{1}\right|=2^{m}-1$ and $\left|E_{m}^{2}\right|=3 \times 2^{m-1}-1$.
(2) Define $\mathcal{E}:=\left\{E_{m}^{i} \mid i=1,2, m \geq 1\right\}$, which is called the envelope set of the Perioddoubling sequence.
(3) We define the order of envelope words that $E_{m}^{1} \sqsubset E_{m}^{2}$ and $E_{m}^{i} \sqsubset E_{m+1}^{j}$ for $i, j \in\{1,2\}$, $m \geq 1$. For any factor $\omega \prec \mathbb{D}$, we define $\operatorname{Env}(\omega)=\min _{\sqsubset}\left\{E_{m}^{i} \mid \omega \prec E_{m}^{i}, i=1,2, m \geq 1\right\}$, which is called the envelope of factor $\omega$.
(4) We define the order $E_{m, p}^{i} \sqsubset E_{n, q}^{j}$ if $E_{m}^{i} \sqsubset E_{n}^{j}$ for $i, j \in\{1,2\}$, $m, n, p, q \geq 1$. For any factor $\omega_{p}$ where $(\omega, p) \in\left(\Omega_{\mathbb{D}}, \mathbb{N}\right)$, we define $\operatorname{Env}\left(\omega_{p}\right)=\min _{\sqsubset}\left\{E_{m, q}^{i} \mid \omega_{p} \prec E_{m, q}^{i}, i=1,2, m, q \geq 1\right\}$, which is called the envelope of factor $\omega_{p}$.

Obviously, the lengths of all envelope words are odd, except $E_{1}^{2}=a a$. Moreover, by the definition of $E_{m}^{i}$ and $A_{m} \delta_{m}^{-1}=B_{m} \delta_{m+1}^{-1}$, we have $E_{m+1}^{1}=E_{m}^{1} \delta_{m} E_{m}^{1}$ and $E_{m+1}^{2}=E_{m}^{1} \delta_{m} E_{m}^{1} \delta_{m} E_{m}^{1}$ for $m \geq 1$. By induction, all envelope words are palindromes [29].

Here are the first few values of $A_{m}, B_{m}, E_{m}^{1}$ and $E_{m}^{2}$ :

- $\left\{A_{m}\right\}_{m \geq 0}=a, a b, a b a a, a b a a a b a b, a b a a a b a b a b a a a b a a, \ldots$.
- $\left\{B_{m}\right\}_{m \geq 0}=b, a a, a b a b, a b a a a b a a, a b a a a b a b a b a a a b a b, \ldots$
- $\left\{E_{m}^{1}\right\}_{m \geq 1}=a, a b a, a b a a a b a, a b a a a b a b a b a a a b a, \ldots$.
- $\left\{E_{m}^{2}\right\}_{m \geq 1}=a a, a b a b a, a b a a a b a a a b a, a b a a a b a b a b a a a b a b a b a a a b a, \ldots$

By this definition, there exists a unique envelope for each factor. For instance, $\operatorname{Env}(a b a b)=$ $a b a b a=E_{2}^{2}$ and $\operatorname{Env}(a b a a a b a b)=E_{4}^{1}$. We can check them by the expressions of $E_{m}^{i}$. More precisely, there exists an integer $j$ and two words $u$ and $\nu$, such that

$$
\begin{equation*}
\omega=\operatorname{Env}(\omega)[j+1, j+|\omega|] \text { and } \operatorname{Env}(\omega)=u \cdot \omega \cdot \nu \tag{2.3}
\end{equation*}
$$

where $0 \leq j \leq|\operatorname{Env}(\omega)|-|\omega|,|u|=j, u \triangleleft \operatorname{Env}(\omega)$ and $\nu \triangleright \operatorname{Env}(\omega)$. We proved that the integer $j$ is unique for any fixed $\omega$ in [29].

### 2.2.2. Derived sequences

In Huang-Wen [29, Theorem 2.6], we determined two types of derived sequences.
Theorem 2.9 Let factor $\omega \prec \mathbb{D}$ have expression in Eq. (2.3).
(1) If there exists an integer $m \geq 1$ such that $\operatorname{Env}(\omega)=E_{m}^{1}$, derived sequence $\mathcal{D}_{\omega}(\mathbb{D})=$ $\mathbb{D}(\alpha, \beta \beta)$. More precisely,

$$
\left\{\begin{array}{lll}
\alpha=\Lambda_{\mathbb{D}, \omega}\left(R_{\mathbb{D}, 1}(\omega)\right), & R_{\mathbb{D}, 1}(\omega)=A_{m}, & \left|R_{\mathbb{D}, 1}(\omega)\right|=2^{m}  \tag{2.4}\\
\beta=\Lambda_{\mathbb{D}, \omega}\left(R_{\mathbb{D}, 2}(\omega)\right), & R_{\mathbb{D}, 2}(\omega)=A_{m-1}, & \left|R_{\mathbb{D}, 2}(\omega)\right|=2^{m-1}
\end{array}\right.
$$

Moreover $R_{\mathbb{D}, 0}(\omega)=\mathbb{D}\left[1,(\omega)_{1}-1\right]=u$ and $\left|R_{\mathbb{D}, 0}(\omega)\right|=j$.
(2) If there exists an integer $m \geq 1$ such that $\operatorname{Env}(\omega)=E_{m}^{2}$, derived sequence $\mathcal{D}_{\omega}(\mathbb{D})=$
$\mathbb{D}(\alpha \beta, \alpha \gamma \alpha \gamma)$. More precisely,

$$
\left\{\begin{array}{lll}
\alpha=\Lambda_{\mathbb{D}, \omega}\left(R_{\mathbb{D}, 1}(\omega)\right), & R_{\mathbb{D}, 1}(\omega)=A_{m-1}, & \left|R_{\mathbb{D}, 1}(\omega)\right|=2^{m-1}  \tag{2.5}\\
\beta=\Lambda_{\mathbb{D}, \omega}\left(R_{\mathbb{D}, 2}(\omega)\right), & R_{\mathbb{D}, 2}(\omega)=A_{m-1} A_{m} B_{m+1}, & \left|R_{\mathbb{D}, 2}(\omega)\right|=7 \times 2^{m-1} \\
\gamma=\Lambda_{\mathbb{D}, \omega}\left(R_{\mathbb{D}, 4}(\omega)\right), & R_{\mathbb{D}, 4}(\omega)=B_{m} B_{m-1}, & \left|R_{\mathbb{D}, 4}(\omega)\right|=3 \times 2^{m-1}
\end{array}\right.
$$

Moreover $R_{\mathbb{D}, 0}(\omega)=\mathbb{D}\left[1,(\omega)_{1}-1\right]=A_{m} u$ and $\left|R_{\mathbb{D}, 0}(\omega)\right|=2^{m}+j$.

| $\mathbb{D}$ | $=$ | $a$ | $b$ | $a$ | $a$ | $a$ | $b$ | $a$ | $b$ | $a$ | $b$ | $a$ | $a$ | $a$ | $b$ | $a$ | $\cdots$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathbb{D}(\alpha, \beta \beta)$ | $=$ | $\alpha$ | $\beta \beta$ | $\alpha$ | $\alpha$ | $\alpha$ | $\beta \beta$ | $\alpha$ | $\beta \beta$ | $\alpha$ | $\beta \beta$ | $\alpha$ | $\alpha$ | $\alpha$ | $\beta \beta$ | $\alpha$ | $\cdots$ |
| $\mathbb{D}(\alpha \beta, \alpha \gamma \alpha \gamma)$ | $=$ | $\alpha \beta \alpha \gamma \alpha \gamma \alpha \beta \alpha \beta \alpha \beta$ | $\alpha \gamma \alpha \gamma$ | $\alpha \beta$ | $\alpha \gamma \alpha \gamma \alpha \beta \alpha \gamma \alpha \gamma \alpha \beta \alpha \beta \alpha \beta \alpha \gamma \alpha \gamma \alpha \beta$ | $\cdots$ |  |  |  |  |  |  |  |  |  |  |  |

Figure 2 The first few letters of $\mathbb{D}, \mathbb{D}(\alpha, \beta \beta)$ and $\mathbb{D}(\alpha \beta, \alpha \gamma \alpha \gamma)$
Example $2.10\left(\right.$ Derived sequence $\left.\mathcal{D}_{\omega}(\mathbb{D})\right)$ Let $\omega=a b a, \operatorname{Env}(a b a)=a b a=E_{2}^{1}$ and $\mathcal{D}_{a b a}(\mathbb{D})=$ $\mathbb{D}(\alpha, \beta \beta)$ where $\Theta_{\mathbb{D}, a b a}(\alpha)=R_{\mathbb{D}, 1}(a b a)=a b a a$ and $\Theta_{\mathbb{D}, a b a}(\beta)=R_{\mathbb{D}, 2}(a b a)=a b$.

$$
\begin{equation*}
\mathbb{D}=\underbrace{a b a a}_{\alpha} \underbrace{a b}_{\beta} \underbrace{a b}_{\beta} \underbrace{a b a a}_{\alpha} \underbrace{a b a a}_{\alpha} \underbrace{a b a a}_{\alpha} \underbrace{a b}_{\beta} \underbrace{a b}_{\beta} \underbrace{a b a a}_{\alpha} \underbrace{a b}_{\beta} \underbrace{a b}_{\beta} \cdots \tag{2.6}
\end{equation*}
$$

Let $\omega=a a, \operatorname{Env}(a a)=a a=E_{1}^{2}$. In this case $\mathbb{D}(\alpha \beta, \alpha \gamma \alpha \gamma)$ where $\Theta_{\mathbb{D}, a a}(\alpha)=R_{\mathbb{D}, 1}(a a)=a$, $\Theta_{\mathbb{D}, a a}(\alpha)=R_{\mathbb{D}, 2}(a a)=a a b a b a b$ and $\Theta_{\mathbb{D}, a a}(\alpha)=R_{\mathbb{D}, 4}(a a)=a a b$.


### 2.2.3. Three steps to prove Theorem 2.9

How to prove Theorem 2.9 by the envelope word? We need three steps:
Step 1. Determine derived sequence $\mathcal{D}_{\omega}(\rho)$ for any envelope word $\omega \prec \mathcal{E}$.
Step 2. For $\omega \prec \rho$, prove that the difference $\operatorname{occ}(\omega, p)-\operatorname{occ}(\operatorname{Env}(\omega), p)$ is independent of $p$.
Step 3. Prove derived sequence $\mathcal{D}_{\omega}(\rho)$ is exactly $\mathcal{D}_{\operatorname{Env}(\omega)}(\rho)$.
For details [29].


Figure 3 The relation between $\omega_{p}$ and $\operatorname{Env}(\omega)_{p}$ for $\omega=b a a$ and $p=1,2,3,4$

### 2.2.4. The reflexivity of derived sequences

A known result in that for any $\omega \prec \rho$ and $v \prec \mathcal{D}_{\omega}(\rho)$, there exists a factor $\mu \prec \rho$ such
that derived sequence $\mathcal{D}_{v}\left(\mathcal{D}_{\omega}(\rho)\right)=\mathcal{D}_{\mu}(\rho)$ (see $[2$, Proposition 6(5)]). Thus for any $\omega \prec \mathbb{D}$ and $v \prec \mathcal{D}_{\omega}(\mathbb{D})$, derived sequence $\mathcal{D}_{v}\left(\mathcal{D}_{\omega}(\mathbb{D})\right) \in\left\{\mathcal{D}_{\omega}(\mathbb{D}) \mid \omega \prec \mathbb{D}\right\}=\{\mathbb{D}(\alpha, \beta \beta), \mathbb{D}(\alpha \beta, \alpha \gamma \alpha \gamma)\}$.

Let $\rho \in\{\mathbb{D}(\alpha, \beta \beta), \mathbb{D}(\alpha \beta, \alpha \gamma \alpha \gamma)\}$. In Huang-Wen [29], we first define the envelope words in $\rho$. Then we divide $\Omega_{\rho}$ into two types: $\Omega_{\rho}^{1}$ and $\Omega_{\rho}^{2}$, according to their envelopes. At last, we prove that for all $v \in \Omega_{\rho}^{1}, \mathcal{D}_{v}(\rho)=\mathbb{D}(\alpha, \beta \beta)$; otherwise $\mathcal{D}_{v}(\rho)=\mathbb{D}(\alpha \beta, \alpha \gamma \alpha \gamma)$. We call it the reflexivity property of derived sequence, see Figure 4.


Figure 4 The reflexivity of derived sequences. For instance, the edge ${ }^{\prime} \mathbb{D} \xrightarrow{\operatorname{Env}(\omega)=E_{m}^{1}} \mathbb{D}(\alpha, \beta \beta) \prime$ means that for any $\omega \prec \mathbb{D}$ if $\operatorname{Env}(\omega)=E_{m}^{1}$ then derived sequence $\mathcal{D}_{\omega}(\mathbb{D})=\mathbb{D}(\alpha, \beta \beta)$, which is given in Theorem 2.9 (1)

## 3. The factor spectrum

Recall the definition of factor spectrum given in the first paragraph in Section 1. Our aim in this section is to determine the factor spectra for some combinatorics properties, such as separated $\left(\mathcal{P}_{1}\right)$, adjacent $\left(\mathcal{P}_{2}\right)$ and overlapped $\left(\mathcal{P}_{3}\right)$. The structure of derived sequences will play an important role in these studies.

For instance, $\omega_{p} \in \mathcal{P}_{2}$ means that there exists an integer $q(>p)$ such that the $p$-th and $q$-th occurrences of $\omega$ are adjacent. And $\omega \in \mathcal{P}_{2}$ means that there exist two integers $p$ and $q$ such that the $p$-th and $q$-th occurrences of $\omega$ are adjacent. In this case, $\omega \omega \prec \rho$. We call $\omega \omega$ a square in $\rho$. We consider $\omega=a b \in \mathbb{D}$ for example.

$$
\begin{equation*}
\mathbb{D}=\underbrace{a b}_{[1]} a a \underbrace{a b}_{[2]} \underbrace{a b}_{[3]} \underbrace{a b}_{[4]} a a \underbrace{a b}_{[5]} a a \cdots . \tag{3.1}
\end{equation*}
$$

We denote the first five occurrences of $\omega$ by notations [1] to [5] in Eq. (3.1). Thus the first five
occurrences of $\omega$ are $\mathbb{D}[1,2], \mathbb{D}[5,6], \mathbb{D}[7,8], \mathbb{D}[9,10]$ and $\mathbb{D}[13,14]$.
$\left\{\begin{aligned} \text { The first (resp., } & 4 \text {-th, } 5 \text {-th) occurrence of } \omega \text { is not followed by another } \omega=a b ; \\ & \Longrightarrow \omega_{1} \notin \mathcal{P}_{2}, \omega_{4} \notin \mathcal{P}_{2}, \omega_{5} \notin \mathcal{P}_{2} ; \\ & \Longrightarrow \text { square } \omega \omega=a b a b \text { does not exist at these positions. } \\ \text { The second } & \text { (resp., third) occurrence of } \omega \text { is followed by another } \omega=a b ; \\ & \Longrightarrow \omega_{2} \in \mathcal{P}_{2}, \omega_{3} \in \mathcal{P}_{2} ; \\ & \Longrightarrow \text { square } \omega \omega=a b a b \text { exists at these positions. }\end{aligned}\right.$

### 3.1. The factor spectra in the Fibonacci sequence

Huang-Wen [3] discussed the factor spectra for some combinatorial properties in the Fibonacci sequence. We first define several subsets of $\left(\Omega_{\mathbb{F}}, \mathbb{N}\right)$ such that the disjoint union of them is $\left(\Omega_{\mathbb{F}}, \mathbb{N}\right)$, see Eq. (3.2) and Figure 5.

$$
\left\{\begin{array}{l}
S_{1}=\bigcup_{m=-1}\left\{(\omega, p)\left|\operatorname{Ker}(\omega)=K_{m},|\omega|=f_{m}\right\},\right.  \tag{3.2}\\
S_{2.1}=\bigcup_{m=-1}\left\{(\omega, p)\left|\operatorname{Ker}(\omega)=K_{m},|\omega|=f_{m+1}, \mathbb{F}[p]=a\right\},\right. \\
S_{2.2}=\bigcup_{m=-1}\left\{(\omega, p)\left|\operatorname{Ker}(\omega)=K_{m},|\omega|=f_{m+1}, \mathbb{F}[p]=b\right\},\right. \\
S_{3.1}=\bigcup_{m=-1}\left\{(\omega, p)\left|\operatorname{Ker}(\omega)=K_{m},|\omega|=f_{m+2}, \mathbb{F}[p]=a\right\},\right. \\
S_{3.2}=\bigcup_{m=-1}\left\{(\omega, p)\left|\operatorname{Ker}(\omega)=K_{m},|\omega|=f_{m+2}, \mathbb{F}[p]=b\right\},\right. \\
S_{4}=\bigcup_{m=-1}\left\{(\omega, p)\left|\operatorname{Ker}(\omega)=K_{m}, f_{m}<|\omega|<f_{m+1}\right\},\right. \\
S_{5.1}=\bigcup_{m=-1}\left\{(\omega, p)\left|\operatorname{Ker}(\omega)=K_{m}, f_{m+1}<|\omega|<f_{m+2}, \mathbb{F}[p]=a\right\},\right. \\
S_{5.2}=\bigcup_{m=-1}\left\{(\omega, p)\left|\operatorname{Ker}(\omega)=K_{m}, f_{m+1}<|\omega|<f_{m+2}, \mathbb{F}[p]=b\right\},\right. \\
S_{6}=\bigcup_{m=-1}\left\{(\omega, p)\left|\operatorname{Ker}(\omega)=K_{m}, f_{m+2}<|\omega|<f_{m+3}\right\} .\right.
\end{array}\right.
$$

Obviously, the disjoint union of $S_{2.1}$ and $S_{2.2}$ is $\bigcup_{m=-1}\left\{(\omega, p)\left|\operatorname{Ker}(\omega)=K_{m},|\omega|=f_{m+1}\right\}\right.$.

|  | $\mathbb{F}[p]=a$ | $\mathbb{F}[p]=b$ |  |
| ---: | :--- | :--- | :---: |
| $\|\omega\|=f_{m}$ | $S_{1}$ | $S_{2.2}$ |  |
| $\|\omega\|=f_{m+1}$ | $S_{2.1}$ | $S_{3.2}$ |  |
| $\|\omega\|=f_{m+2}$ | $S_{3.1}$ |  |  |
| $f_{m}<\|\omega\|<f_{m+1}$ | $S_{4}$ | $S_{5.2}$ |  |
| $f_{m+1}<\|\omega\|<f_{m+2}$ | $S_{5.1}$ |  |  |
| $f_{m+2}<\|\omega\|<f_{m+3}$ | $S_{6}$ |  |  |

Figure 5 Several subsets of $\left(\Omega_{\mathbb{F}}, \mathbb{N}\right)$ in Eq. (3.2)

Property $3.1 \operatorname{Spt}\left(\mathbb{F}, \mathcal{P}_{1}\right)=S_{1} \cup S_{2.1} \cup S_{4} \cup S_{5.1}, \quad \operatorname{Spt}\left(\mathbb{F}, \mathcal{P}_{2}\right)=S_{2.2} \cup S_{3.1}, \quad \operatorname{Spt}\left(\mathbb{F}, \mathcal{P}_{3}\right)=$ $S_{3.2} \cup S_{5.2} \cup S_{6}$.

Using the factor spectrum $\operatorname{Spt}\left(\mathbb{F}, \mathcal{P}_{2}\right)$, we get a new proof of the conclusion that all squares in $\mathbb{F}$ are of length $2 f_{m}$. Furthermore, we find that not all occurrences of all factors of length $f_{m}$ belong to $\operatorname{Spt}\left(\mathbb{F}, \mathcal{P}_{2}\right)$.

### 3.2. The factor spectra in the Period-doubling sequence

Huang-Wen [29] discussed the factor spectra for some combinatorial properties in the Perioddoubling sequence. We first define several subsets of $\left(\Omega_{\mathbb{D}}, \mathbb{N}\right)$ such that the disjoint union of them is $\left(\Omega_{\mathbb{D}}, \mathbb{N}\right)$, see Eq. (3.3) and Figure 6.

$$
\left\{\begin{array}{l}
S_{1}=\bigcup_{m=1}\left\{(\omega, p) \mid \operatorname{Env}(\omega)=E_{m}^{1}, \mathbb{D}(\alpha, \beta \beta)[p]=\alpha\right\},  \tag{3.3}\\
S_{2.1}=\bigcup_{m=1}\left\{(\omega, p)\left|\operatorname{Env}(\omega)=E_{m}^{1},|\omega|<2^{m-1}, \mathbb{D}(\alpha, \beta \beta)[p]=\beta\right\},\right. \\
S_{2.2}=\bigcup_{m=1}\left\{(\omega, p)\left|\operatorname{Env}(\omega)=E_{m}^{1},|\omega|=2^{m-1}, \mathbb{D}(\alpha, \beta \beta)[p]=\beta\right\},\right. \\
S_{2.3}=\bigcup_{m=1}\left\{(\omega, p)\left|\operatorname{Env}(\omega)=E_{m}^{1},|\omega|>2^{m-1}, \mathbb{D}(\alpha, \beta \beta)[p]=\beta\right\},\right. \\
S_{3}=\bigcup_{m=1}\left\{(\omega, p) \mid \operatorname{Env}(\omega)=E_{m}^{2}, \mathbb{D}(\alpha \beta, \alpha \gamma \alpha \gamma)[p]=\alpha\right\}, \\
S_{4}=\bigcup_{m=1}\left\{(\omega, p) \mid \operatorname{Env}(\omega)=E_{m}^{2}, \mathbb{D}(\alpha \beta, \alpha \gamma \alpha \gamma)[p] \neq \alpha\right\} .
\end{array}\right.
$$

The disjoint union of $S_{2.1}, S_{2.2}$ and $S_{2.3}$ is $\bigcup_{m=1}\left\{(\omega, p) \mid \operatorname{Env}(\omega)=E_{m}^{1}, \mathbb{D}(\alpha, \beta \beta)[p]=\beta\right\}$.

|  | $\mathbb{D}(\alpha, \beta \beta)[p]=\alpha$ | $\mathbb{D}(\alpha, \beta \beta)[p]=\beta$ |
| :--- | :---: | :---: |
| $\operatorname{Env}(\omega)=E_{m}^{1}$ | $S_{1}$ | $S_{2.1}$ |
|  |  | $S_{2.2}$ |
|  | $S_{3}$ | $S_{2.3}$ |
|  | $\mathbb{D}(\alpha \beta, \alpha \gamma \alpha \gamma)[p]=\alpha$ | $\mathbb{D}(\alpha \beta, \alpha \gamma \alpha \gamma)[p] \neq \alpha$ |

Figure 6 Several subsets of $\left(\Omega_{\mathbb{D}}, \mathbb{N}\right)$ in Eq. (3.3)

Property 3.2 $\operatorname{Spt}\left(\mathbb{D}, \mathcal{P}_{1}\right)=S_{1} \cup S_{4} \cup S_{2.1}, \operatorname{Spt}\left(\mathbb{D}, \mathcal{P}_{2}\right)=S_{2.2}, \operatorname{Spt}\left(\mathbb{D}, \mathcal{P}_{3}\right)=S_{3} \cup S_{2.3}$.
Using $S p t\left(\mathbb{D}, \mathcal{P}_{2}\right)$, we get new proofs of some conclusions in Damanik [13]:
(1) All squares in $\mathbb{D}$ are $\bigcup_{m \geq 1}\left\{A_{m-1}\left[j+1,2^{m-1}\right] A_{m-1} A_{m-1}[1, j] \mid 0 \leq j<2^{m-1}\right\}$;
(2) There are $2^{m-1}$ 's distinct squares with length $2^{m}$ for $m \geq 1$.

## 4. Enumerations

Using Theorems 2.6 and 2.9, we determined the numbers of palindromes, squares and cubes occurring in each factor of $\mathbb{F}, \mathbb{T}$ and $\mathbb{D}$, see Huang-Wen $[11,12,30]$ for instance. These topics are of great importance in computer science.

### 4.1. Enumeration of palindromes in $\mathbb{F}$ and $\mathbb{T}$

A palindrome is a finite word that reads the same backwards as forwards. Let Pal ${ }_{\mathbb{F}}$ (resp., $\mathrm{Pal}_{\mathbb{T}}$ ) be all palindromes occurring in $\mathbb{F}$ (resp., $\mathbb{T}$ ). Some previous research has been done on "rich word", which is based on the number of distinct palindromes. A finite word $\omega$ is rich if and only if $\omega$ contains exactly $|\omega|+1$ distinct palindromes (including the empty word). An infinite word is rich if and only if all of its factors are rich. Droubay-Justin-Pirillo [31] proved that episturmian
sequences are rich. Therefore, as special cases, $\mathbb{F}$ and $\mathbb{T}$ are rich. Thus the number of distinct palindromes in $\mathbb{F}[1, n]$ (resp., $\mathbb{T}[1, n]$ ) is $n+1$ for all $n$.

In Huang-Wen [11], we consider the numbers of repeated palindromes in $\mathbb{F}[1, n]$ and $\mathbb{T}[1, n]$. All results in this subsection are from this paper. Denote

$$
A(n)=\#\left\{(\omega, p) \mid \omega \in \operatorname{Pal}_{\mathbb{F}}, \omega_{p} \prec \mathbb{F}[1, n]\right\} \text { and } B(n)=\#\left\{(\omega, p) \mid \omega \in \operatorname{Pal}_{\mathbb{T}}, \omega_{p} \prec \mathbb{T}[1, n]\right\}
$$

The research on counting the repeated palindromes is not rich. From our knowledge, it seems the first work to study this problem. In related fields, the numbers of special types of factors have been investigated in recent years, such as squares, cubes, $r$-powers, palindromes, runs, Lyndon factors, etc $[10,21,31-41]$.

The main difficulty of this problem is twofold:
(1) The positions of all occurrences for all palindromes are not easy to be determined. In [11], we overcome this difficulty by using the derived sequence properties of $\mathbb{F}$ and $\mathbb{T}$, which we introduced and studied in $[3,4]$, see also Theorem 2.6.
(2) Taking $\mathbb{F}$ for instance, by the derived sequence property of $\mathbb{F}$, we can find out all distinct palindromes in $\mathbb{F}[1, n]$. We can also count the number of occurrences of each palindrome. So the summation of these numbers are the numbers of repeated palindromes in $\mathbb{F}[1, n]$. But this method is complicated. We overcome this difficulty by studying the relations among positions of each $\omega_{p}$, and establishing the recursive structure of $\mathrm{Pal}_{\mathbb{F}}$. Using the derived sequence properties and recursive structures, we give algorithms for counting $A(n)$ and $B(n)$, respectively.

Take $A(n)$ for instance. We have $A(n)=\sum_{i=1}^{n} a(i)$ where $a(i)$ is given in Property 4.1.
Property 4.1 The vectors $[a(1)]=[1],[a(2), a(3)]=[1,2]$ and for $m \geq 3$

$$
\begin{aligned}
& {\left[a\left(f_{m}-1\right), \ldots, a\left(f_{m+1}-2\right)\right]} \\
& \quad=\left[a\left(f_{m-2}-1\right), \ldots, a\left(f_{m-1}-2\right), a\left(f_{m-1}-1\right), \ldots, a\left(f_{m}-2\right)\right]+[\underbrace{1, \ldots, 1}_{f_{m-1}}] .
\end{aligned}
$$

The first few values of $a(n)$ are $[a(1)]=[1],[a(2), a(3)]=[1,2]$,

$$
\begin{aligned}
& {[a(4), a(5), a(6)]=[a(1), a(2), a(3)]+[1,1,1]=[2,2,3]} \\
& {[a(7), \ldots, a(11)]=[a(2), \ldots, a(6)]+[1,1,1,1,1]=[2,3,3,3,4]} \\
& {[a(12), \ldots, a(19)]=[a(4), \ldots, a(11)]+[\underbrace{1, \ldots, 1}_{8}]=[3,3,4,3,4,4,4,5] .} \\
& {[a(20), \ldots, a(32)]=[a(7), \ldots, a(19)]+[\underbrace{1, \ldots, 1}_{13}]=[3,4,4,4,5,4,4,5,4,5,5,5,6] .}
\end{aligned}
$$

We also get explicit expressions for some special $n$, such as: for $m \geq 0$,

$$
\left\{\begin{array}{l}
A\left(f_{m}\right)=\frac{m-3}{5} f_{m+2}+\frac{m-1}{5} f_{m}+m+3  \tag{4.1}\\
B\left(t_{m}\right)=\frac{m}{22}\left(10 t_{m}+5 t_{m-1}+3 t_{m-2}\right)+\frac{1}{22}\left(-23 t_{m}+12 t_{m-1}-5 t_{m-2}\right)+m+\frac{3}{2}
\end{array}\right.
$$

We think this method for counting the repeated palindromes is valid for the $m$-bonacci word, and even valid for sturmian sequences, episturmian sequences etc. But now we only have the derived sequence properties of $\mathbb{F}, \mathbb{T}$ and $\mathbb{D}$. In the final remark in [11], we establish the cylinder structures and chain structures of $\mathrm{Pal}_{\mathbb{F}}$ and $\mathrm{Pal}_{\mathbb{T}}$. Using them, we prove some known results.

### 4.2. Enumeration of repetitions in $\mathbb{F}$

The fractional power is a topic dealing with repetitions in words. We say a (finite or infinite) word $\omega$ contains a $r$-power (real $r>1$ ) if $\omega$ has a factor of the form $x^{\lfloor r\rfloor} x^{\prime}$ where $x^{\prime}$ is a prefix of $x$ and $\left|x^{[r]} x^{\prime}\right| \geq r|x|$ see [1]. In this case, we call $x^{\lfloor r\rfloor} x^{\prime}$ a $r$-power with size $|x|$. For instance, taking $x=a b$, then $\mathbb{F}[4,8]=a b a b a$ is a $\frac{5}{2}$-power of size $|a b|=2$ in $\mathbb{F}$. Obviously the notion $r$-power is a generalization of square ( 2 -power) and cube (3-power).

The study of power of a word has a long history. There are many significant contributions, for example $[9,10,32-35,39,42-44]$. In particular, Iliopoulos-Moore-Smyth [35] computed the positions of all squares in $\mathbb{F}$. Fraenkel-Simpson $[32,33]$ obtained the number of squares in $F_{m}=$ $\sigma^{m}(a)$. Damanik-Lenz [43] studied the index of Sturmian sequences. Glen [34] determined all of the squares (and subsequently higher powers) occurring in episturmian words. Using the results in [34], it was possible to determine the exact number of distinct squares in each building block (for instance, the building blocks in the Fibonacci sequence is $F_{m}=\sigma^{m}(a)$ ), which extends Fraenkel and Simpson's result [32,33]. Du-Mousavi-Schaeffer-Shallit [44,45] obtained the numbers of repeated squares and cubes in $\mathbb{F}[1, n]$ for all $n \geq 1$. All their numerations start from the first letter of the sequence.

In Huang-Wen [12], we count the number of distinct $r$-powers in $\mathbb{F}[i, n+i-1]$ for all $i, n \geq 1$ and $r \geq 2$, denoted by $\mathrm{D}(r, i, n)$. Our numeration can start from any letter of the sequence, comparing starting from the first letter, there are some difficulties. To overcome these difficulties, we introduce a new notion called the position sequence. In Huang-Wen [30], we count the number of repeated $r$-powers in $\mathbb{F}[i, n+i-1]$ for all $i, n \geq 1$ and $r \geq 2$, denoted by $\mathrm{R}(r, i, n)$. We give precise results for $r \in\{2,2+\epsilon, 3\}$, where $\epsilon$ is a small positive number.

The methods of counting the distinct or repeated factors in $\mathbb{F}[i, n+i-1]$ is quite different. Take $r=2$ (squares) for instance. Let

$$
\left.h=\left\lfloor\frac{\ln (\sqrt{5} n / 2)}{\ln \alpha}+\frac{2}{n}\right\rfloor-2 \quad \text { (i.e., } 2 f_{h} \leq n<2 f_{h+1}\right) .
$$

(1) A well known result is that the number of distinct squares of length $2 f_{m}$ in the Fibonacci sequence is $f_{m}$ (see [9]). By the arguments in [12, Subsection 5.2], $\mathbb{F}[i, n+i-1]$ contains all distinct squares of length $2 f_{m}$ for any $m \geq 0$ and $n \geq f_{m+3}+2 f_{m}-1$. Thus in order to count the number of distinct square in $\mathbb{F}[i, n+i-1]$ for $n \geq 6$, we only need to pay close attention to the number of the squares of length $f_{m}$ where $m \in\{h-2, h-1, h\}$. Through careful observation and analysis, we get explicit expression of the number of distinct fractional powers in each factor of the Fibonacci sequence.
(2) Obviously, this method above is not fit for counting the repeated squares. Comparing with the distinct case, we will introduce two completely different tools, which we call square tree and matrix decomposition. The former is used for counting the number of repeated $r$-powers in $\mathbb{F}[1, n+i-1]$, i.e., $\mathrm{R}(r, 1, n)$. The latter is used for determining the difference between $\mathrm{R}(r, i, n)$ and $\mathrm{R}(r, 1, n)$.

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